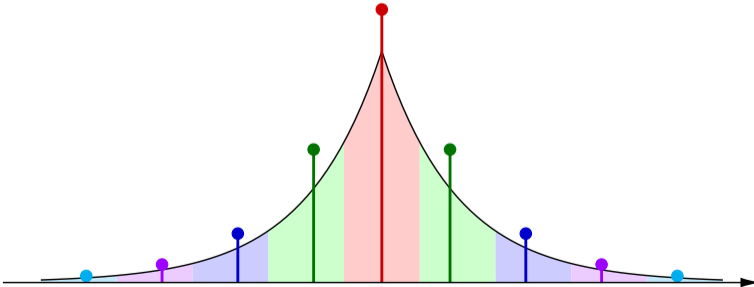


# Scalar Quantization



# Last Lectures: Lossless Coding

## Variable-length Coding

- Scalar codes, conditional codes, block codes, V2V codes (using codeword tables)
- For given pmf: **Huffman algorithm** yields optimal codeword table
- Problem: Codeword tables become too large for practical application of block codes

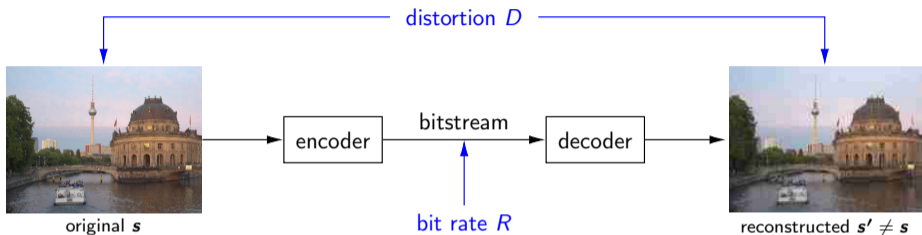
## Arithmetic Coding

- No codeword table: On-the-fly encoding and decoding
- Sub-optimal block code for arbitrarily large block sizes  $N$  (very close to optimum for  $N \gg 1$ )
- Straightforward combination with **conditional and adaptive probability models**

## Reduction of Inter-Symbols Dependencies before Entropy Coding

- **Affine and linear prediction**: Suitable for reducing dependencies in audio, image, video data
- **Lempel-Ziv coding** or **block sorting**: Suitable for text, source code, general files
- **Lossless coding in practice**:
  - Prediction followed by entropy coding of prediction errors
  - Lempel-Ziv coding or block sorting followed by entropy coding

# Lossy Coding

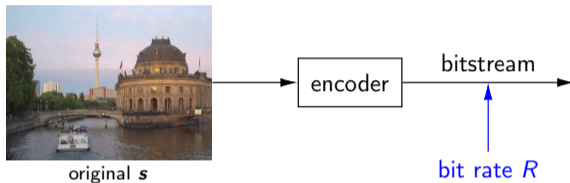


Lossy coding is characterized by two aspects:

- Bit rate  $R$ : Average number of bits per sample (or per time unit)
- Distortion  $D$ : Measure for deviation between original signal  $s$  and reconstructed signal  $s'$

**Design Goal:** Smallest possible bit rate for given maximum distortion, or  
Smallest possible distortion for given maximum bit rate

# Lossy Coding: Bit Rate



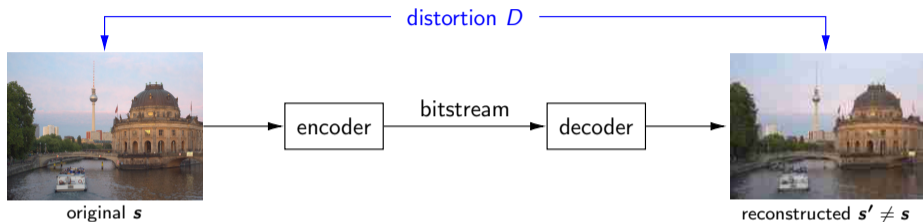
## Bit Rate $R$ :

- Images: Average number of bits per sample
- Audio or video: Average number of bits per time units

## Often used Approximation:

- Assume that we have a close to optimal entropy coding (e.g., arithmetic coding)
- Bit rate = Entropy of symbols that are actually transmitted

# Lossy Coding: Distortion



## Distortion Measures used in Practice

- General  $p$ -norm distortion:

$$D_p = \frac{1}{N} \sum_{k=1}^N |s_k - s'_k|^p \quad \text{or} \quad D_p = \mathbb{E} \left\{ |S - S'|^p \right\}$$

- Most often: Mean squared error (MSE)

$$D_2 = \frac{1}{N} \sum_{k=1}^N (s_k - s'_k)^2 \quad \text{or} \quad D_2 = \mathbb{E} \left\{ (S - S')^2 \right\}$$

# MSE Distortion as Signal-To-Noise Ratio (SNR)

## Signal-to-Noise Ratio (SNR)

- Logarithmic ratio of variance and MSE distortion

$$\text{SNR} = 10 \cdot \log_{10} \left( \frac{\sigma^2}{D_2} \right)$$

- Measured in decibel (dB)

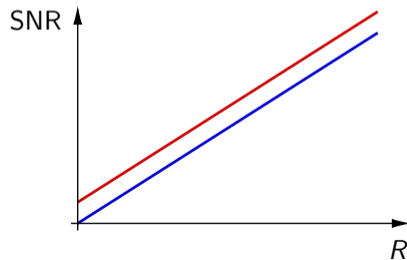
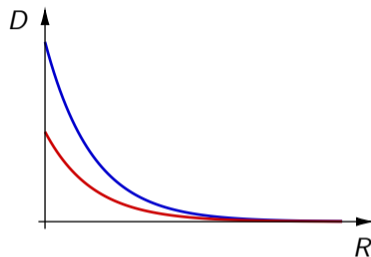
## Advantages of using SNR

- Independent of signal variance
- Easy interpretation of differences

$$\Delta \text{SNR} = \text{SNR}_a - \text{SNR}_b = -10 \cdot \log_{10} \frac{D_a}{D_b}$$

- Examples:
 

$D_a = D_b$	$\rightarrow$	$\Delta \text{SNR} \approx 0.0 \text{ dB}$
$D_a = D_b / \sqrt{2}$	$\rightarrow$	$\Delta \text{SNR} \approx 1.5 \text{ dB}$
$D_a = D_b / 2$	$\rightarrow$	$\Delta \text{SNR} \approx 3.0 \text{ dB}$
$D_a = D_b / 4$	$\rightarrow$	$\Delta \text{SNR} \approx 6.0 \text{ dB}$
$D_a = D_b / 8$	$\rightarrow$	$\Delta \text{SNR} \approx 9.0 \text{ dB}$



# Probabilistic Modeling of Sources

## Source Coding in Practice

- Encoder and decoder are computer programs
- Actual input signals are **discrete-time** and **discrete-amplitude** signals

## Real-world signals

- In most cases: Continuous-time and continuous-amplitude signals
- Discrete signals are obtained by sampling and quantization
- Typical scenarios: Initial quantization has negligible effect on source coding

## Theoretical Analysis of Lossy Source Coding

- Will mostly use models for **discrete-time** and **continuous-amplitude** signals
- Main reason: Mathematical tractability
- Interpretation: Consider signal before initial quantization

# Review: Random Variables and Cumulative Distribution Function (CDF)

## Random Variable

- Function  $X(\zeta)$  of the sample space  $\mathcal{O}$  that assigns a real value  $x = X(\zeta)$  to each possible outcome  $\zeta \in \mathcal{O}$  of a random experiment

## Cumulative Distribution Function (cdf)

- Cumulative distribution function  $F_X(x)$  of a random variable  $X$

$$F_X(x) = P(X \leq x) = P(\{\zeta : X(\zeta) \leq x\})$$

- Joint cdf of two random variables  $X$  and  $Y$

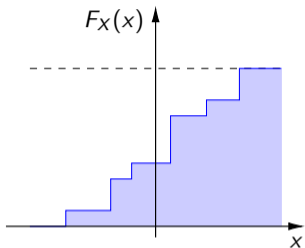
$$F_{XY}(x, y) = P(X \leq x, Y \leq y)$$

- Conditional cdf of a random variable  $X$  given another random variable  $Y$

$$F_{X|Y}(x|y) = P(X \leq x | Y \leq y) = \frac{P(X \leq x, Y \leq y)}{P(Y \leq y)} = \frac{F_{XY}(x, y)}{F_Y(y)}$$



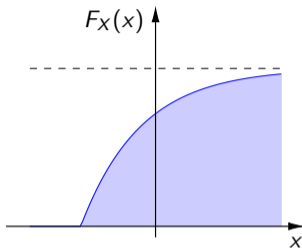
# Review: Examples of Cumulative Distribution Functions



## Staircase function

- Random variable  $X$  can only take a countable number of values

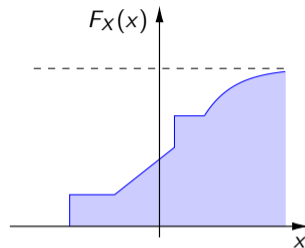
→ **Discrete random variable**



## Continuous function

- Random variable  $X$  can take all values inside one or more non-zero intervals

→ **Continuous random variable**



## Mixed type

- Random variable  $X$  can take all values inside one or more non-zero intervals and a countable number of additional values

# Continuous Random Variables and Probability Density Function (PDF)

## Continuous Random Variables

- A random variable  $X$  is called a **continuous random variable** if and only if its cdf  $F_X(x)$  is a continuous function

## Probability Density Function

- Probability density function (pdf) of a continuous random variable  $S$

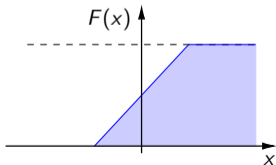
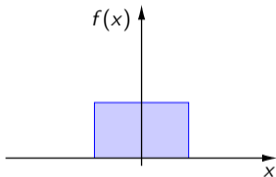
$$f_X(x) = \frac{\partial}{\partial x} F_X(x) \iff F_X(x) = \int_{-\infty}^x f_X(t) dt$$

- Properties:
- $f_X(x) \geq 0, \forall x$
  - $\int_{-\infty}^{\infty} f_X(t) dt = 1$
  - $P(a < X \leq b) = \int_a^b f_X(t) dt$

## Examples for Continuous Distributions (Zero Mean)

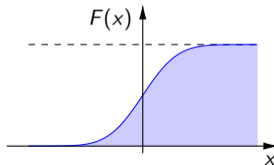
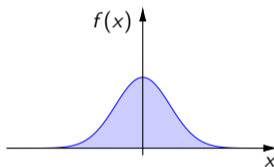
## Uniform

$$f(x) = \begin{cases} \frac{1}{2a} & : |x| \leq a \\ 0 & : \text{otherwise} \end{cases}$$



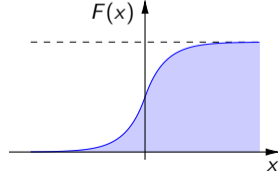
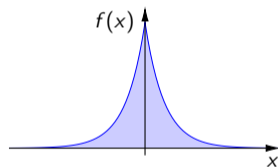
## Gaussian

$$f(x) = \frac{1}{\sqrt{2\pi\sigma^2}} e^{-\frac{(x-\mu)^2}{2\sigma^2}}$$



## Laplacian

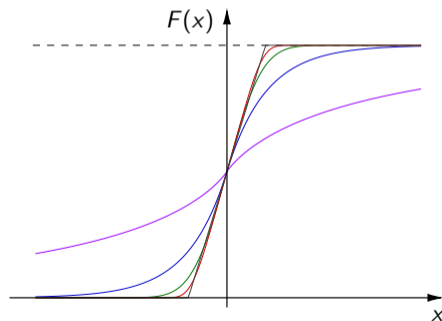
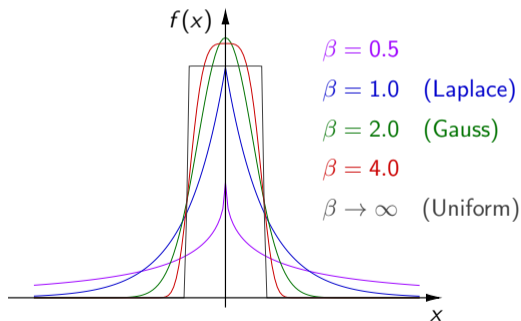
$$f(x) = \frac{1}{\sqrt{2\sigma^2}} e^{-\sqrt{\frac{2}{\sigma^2}} |x-\mu|}$$



# Generalized Gaussian Distribution

Shape parameter  $\beta \in (0, \infty)$ :

$$f(x) = \frac{\beta}{2\alpha\Gamma(1/\beta)} e^{-\left(\frac{|x-\mu|}{\alpha}\right)^\beta} \quad \text{with} \quad \Gamma(x) = \int_0^\infty e^{-t} t^{x-1} dt$$



**→ Suitable approximation for many distributions**

# Joint and Conditional Probability Density Function

## Joint Probability Density Function

- Joint pdf of two random variables  $X$  and  $Y$

$$f_{XY}(x, y) = \frac{\partial^2}{\partial x \partial y} F_{XY}(x, y)$$

## Conditional Probability Density Function

- Conditional pdf of a random variable  $X$  given another random variable  $Y$

$$f_{X|Y}(x|y) = \frac{\partial}{\partial x} F_{X|Y}(x|y) = \frac{\partial}{\partial x} \frac{F_{XY}(x, y)}{F_Y(y)} = \frac{\frac{\partial^2}{\partial x \partial y} F_{XY}(x, y)}{\frac{\partial}{\partial y} F_Y(y)} = \frac{f_{XY}(x, y)}{f_Y(y)}$$

# Expected Values for Continuous Random Variables

## Expected Values

- Expected value of a function  $g(X)$  of a continuous random variable  $X$

$$E\{g(X)\} = E_X\{g(X)\} = \int_{-\infty}^{\infty} g(x) f_X(x) dx$$

- Expected value of function  $g(X, Y)$  of two continuous random variables  $X$  and  $Y$

$$E\{g(X, Y)\} = E_{XY}\{g(X, Y)\} = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} g(x, y) f_{XY}(x, y) dx dy$$

## Conditional Expected Values

- Expected value of function  $g(X)$  of a random variable  $X$  given another random variable  $Y$

$$E\{g(X) | Y\} = \int_{-\infty}^{\infty} g(x) f_{X|Y}(x | Y) dx \quad (\text{is another random variable})$$

# Properties of Expected Values

*same properties as in discrete case*

## Important Properties

- Linearity of expected values

$$E\{ aX + bY \} = a \cdot E\{ X \} + b \cdot E\{ Y \}$$

- For independent random variables  $X$  and  $Y$

$$E\{ XY \} = E\{ X \} E\{ Y \}$$

- Iterative expectation rule

$$E\{ E\{ g(X) | Y \} \} = E\{ g(X) \}$$

## Important Expected Values

- **Mean**  $\mu_X$  of a random variable  $X$

$$\mu_X = E\{X\} = \int_{-\infty}^{\infty} x \cdot f_X(x) dx$$

- **Variance**  $\sigma_X^2$  of a random variable  $X$

$$\sigma_X^2 = E\left\{ (X - E\{X\})^2 \right\} = \int_{-\infty}^{\infty} (x - \mu_X)^2 \cdot f_X(x) dx$$

- **Covariance**  $\sigma_{XY}^2$  of two random variables  $X$  and  $Y$ , and **correlation coefficient**  $\phi_{XY}$

$$\sigma_{XY}^2 = E\left\{ (X - E\{X\}) (Y - E\{Y\}) \right\} = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} (x - \mu_X)(y - \mu_Y) \cdot f_{XY}(x, y) dx dy$$

$$\phi_{XY} = \frac{\sigma_{XY}^2}{\sqrt{\sigma_X^2 \cdot \sigma_Y^2}}$$



# Continuous Random Processes

## Discrete-Time Random Process

- Series of random experiments at time instants  $t_n$ , with  $n = 0, 1, 2, \dots$
- For each experiment: Random variable  $X_n = X(t_n)$
- **Random process**: Series of random variables

$$\mathbf{X} = \{X_0, X_1, X_2, \dots\} = \{X_n\}$$

## Discrete-Time Continuous-Amplitude Random Process

- Random variables  $X_n$  are continuous random variables
- Type of random processes we consider for analyzing lossy coding

# Statistical Properties of Continuous Random Processes

## Characterization of Statistical Properties

- Consider  $N$ -dimensional random vector

$$\mathbf{X}_k^{(N)} = \{X_k, X_{k+1}, \dots, X_{k+N-1}\}$$

- $N$ -th order joint cdf

$$F_k^{(N)}(\mathbf{x}) = \mathrm{P}\left(\mathbf{X}_k^{(N)} \leq \mathbf{x}\right) = \mathrm{P}(X_k \leq x_0, X_{k+1} \leq x_1, \dots, X_{k+N-1} \leq x_{N-1})$$

- $N$ -th order joint pdf

$$f_k^{(N)}(\mathbf{x}) = \frac{\partial^N}{\partial x_0 \cdots \partial x_{N-1}} F_k^{(N)}(\mathbf{x})$$

- Also: Conditional cdfs and conditional pdfs

# Models for Random Processes

## Stationary Random Processes

- Statistical properties are invariant to a shift in time
- In this course: Typically restrict our considerations to stationary processes

## Memoryless Random Processes

- All random variables  $X_n$  are independent of each other

## Independent and Identically Distributed (IID) Random Processes

- Random processes that are **stationary** and **memoryless**

## Markov Processes

- Markov property: Future outcomes do only depend on present outcome, but not on past outcomes

$$F(x_n | x_{n-1}, x_{n-2}, x_{n-3}, \dots) = F(x_n | x_{n-1})$$

$$f(x_n | x_{n-1}, x_{n-2}, x_{n-3}, \dots) = f(x_n | x_{n-1})$$

- Simple model for random processes with memory

# Autoregressive (AR) Processes

## General AR(p) Model

- Autoregressive model of order  $p$  for random variables  $X_n$  with mean  $\mu$

$$X_n = Z_n + \mu + \sum_{k=1}^p \varrho_k \cdot (X_{n-k} - \mu)$$

where  $\mathbf{Z} = \{Z_n\}$  is a zero-mean iid process (innovation process)  
and  $\varrho_1, \dots, \varrho_p$  are the model parameters

## Special case: AR(1) model

- Autoregressive model of order  $p = 1$

$$X_n = Z_n + \mu + \varrho \cdot (X_{n-1} - \mu)$$

- Completely specified by mean  $\mu$ , correlation coefficient  $\varrho$ , and pdf  $f_Z(z)$  of iid process  $\{Z_n\}$
- Important type of stationary Markov process for continuous random processes

# Gaussian Processes

## Gaussian Random Process

- All finite collections of random variables  $X_n$  are Gaussian random vectors
- $N$ -th order pdf is given by  $N$ -th order auto-covariance matrix  $\mathbf{C}_N$  and mean  $\mu$

$$f_{\mathbf{X}}(\mathbf{x}) = \frac{1}{\sqrt{(2\pi)^N |\mathbf{C}_N|}} e^{-\frac{1}{2}(\mathbf{x}-\mu)^T \mathbf{C}_N^{-1}(\mathbf{x}-\mu)} \quad \text{with} \quad \mu = \begin{pmatrix} \mu \\ \vdots \\ \mu \end{pmatrix}$$

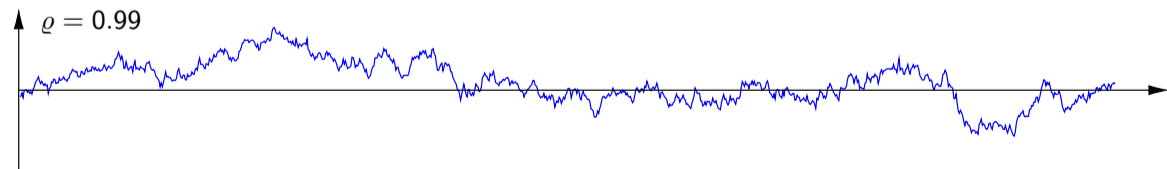
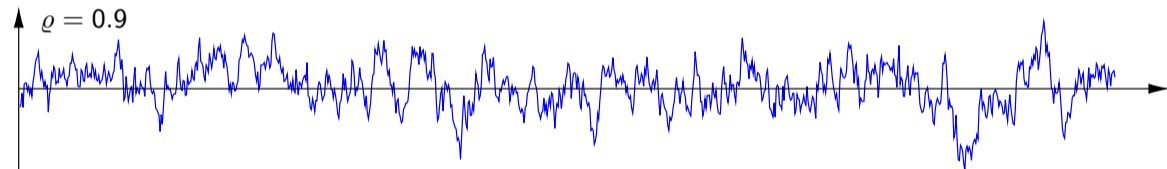
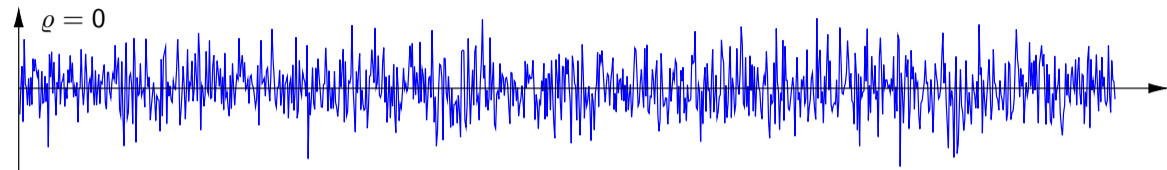
## Stationary Gauss-Markov Process

- Stationary Markov process that is also a Gaussian random process
- Can be constructed with Gaussian iid process  $\mathbf{Z} = \{Z_n\}$  according to

$$X_n = \mu + \varrho(X_{n-1} - \mu) + Z_n$$

- Statistical properties are completely specified by **mean**  $\mu$ , **variance**  $\sigma^2$ , **correlation coefficient**  $\varrho$
- ➔ Will use it as very simple model for sources with memory

# Examples of Gauss-Markov Processes (1000 Samples)



# Summary of Mathematical Basics (for Continuous Case)

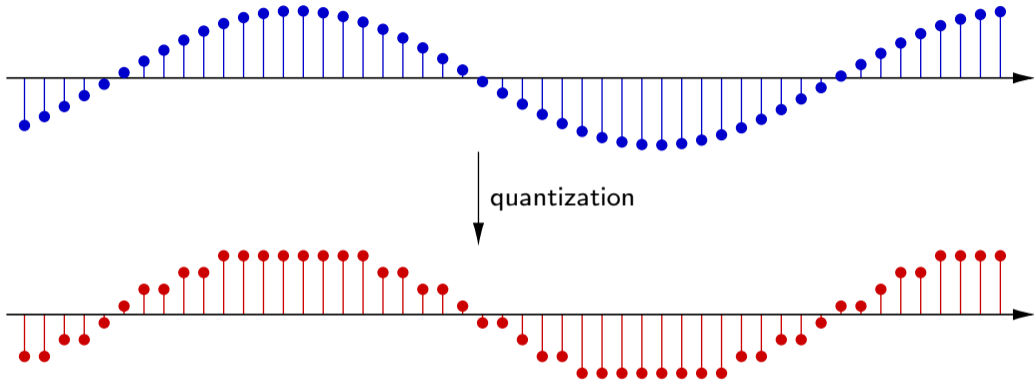
## Continuous Random Variables

- Can take all values inside one or more non-zero intervals
- Cumulative distribution function (cdf): Continuous function
- Probability density function (pdf)
- Expected values: Mean, variance, covariance

## Discrete-Time Continuous-Amplitude Random Processes

- Sequence of continuous random variables: Model for lossy source coding
- Types of random processes: Stationary, memoryless, iid, Markov
- Suitable model for real signals: Autoregressive processes
- Special importance for lossy source coding: Gaussian processes
- Simple model for sources with memory: Gauss-Markov process

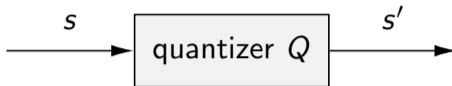
# Quantization



- “Lossy part” of source coding
- Non-reversible mapping from input to output samples
- Determines trade-off between signal fidelity and bit rate



# Scalar Quantization: Functional Mapping



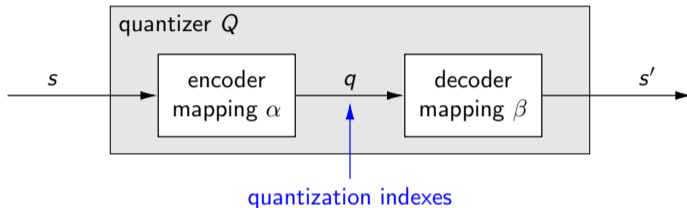
- Scalar Quantization: Functional mapping of an input sample to an output sample

$$s' = Q(s)$$

- Input:
  - Discrete or continuous
- Output:
  - Set of obtainable output points is countable
  - Less obtainable output points than input points

→ **Non-reversible loss in signal fidelity**

# Structure of Scalar Quantizers: Encoder and Decoder Mapping



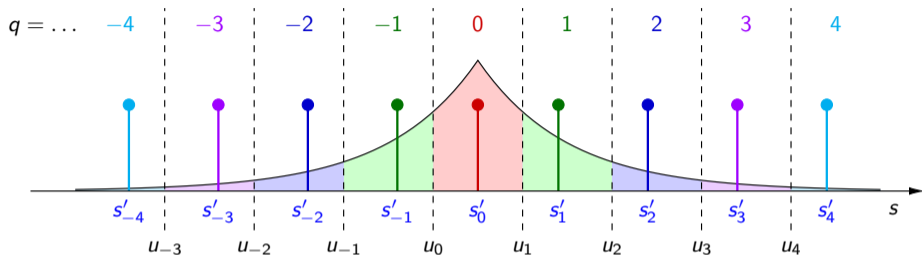
- Split quantizer  $Q$  into **encoder mapping**  $\alpha$  and **decoder mapping**  $\beta$
- Encoder mapping  $\alpha$ : Maps input sample  $s$  to a quantizer index  $q$  (integer)

$$q = \alpha(s)$$

- Decoder mapping  $\beta$ : Maps quantizer index  $q$  to reconstructed samples  $s'$

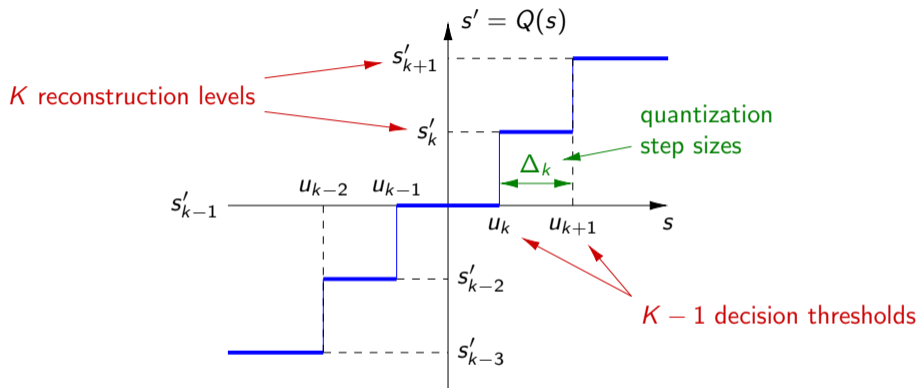
$$s' = \beta(q) = \beta(\alpha(s)) = Q(s)$$

# Principle of Scalar Quantization



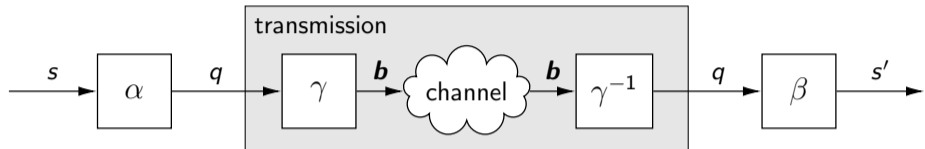
- Partition real line into a countable (typically finite) number of quantization intervals  $\mathcal{I}_k$ 
    - Partitioning is given by decision thresholds  $\{u_k\}$
    - Quantization intervals are labeled by quantization index  $q$
    - A quantization interval is given by  $\mathcal{I}_k = [u_k, u_{k+1})$
  - Each quantization interval  $\mathcal{I}_k$  is associated with a reconstruction level  $s'_k \in \mathcal{I}_k$
- **Scalar quantization:** Replace input value  $s$  that falls inside  $\mathcal{I}_k$  with reconstruction value  $s'_k$

# Scalar Quantization: Input-Output Function



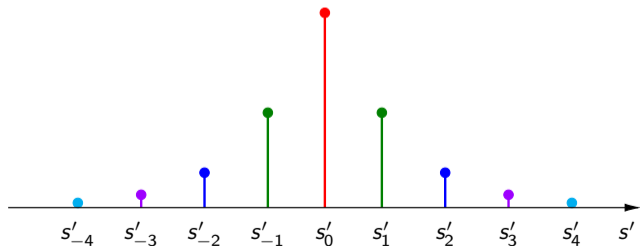
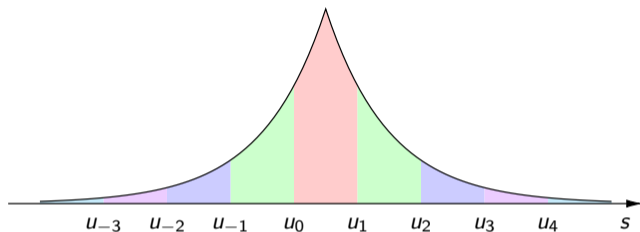
- Scalar quantizer mapping:  $Q : \mathbb{R} \mapsto \{\dots, s'_{k-1}, s'_k, s'_{k+1}, \dots\}$
- Quantization intervals:  $\mathcal{I}_k = [u_k, u_{k+1})$
- Quantization step sizes:  $\Delta_k = u_{k+1} - u_k$

# Scalar Quantization and Entropy Coding



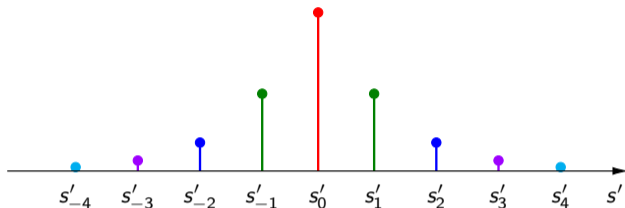
- Add **lossless coding**  $\gamma$  of quantization indexes (e.g., Huffman or arithmetic coding)
- Encoding/decoding process:
  - ① Encoder mapping  $\alpha$ : Input samples  $s \mapsto$  quantization indexes  $q$
  - ② Lossless mapping  $\gamma$ : Quantization indexes  $q \mapsto$  bitstream  $b$
  - ③ Transmission channel: Transmission of bitstream (assume: error-free)
  - ④ Lossless mapping  $\gamma^{-1}$ : Bitstream  $b \mapsto$  quantization indexes  $q$
  - ⑤ Decoder mapping  $\beta$ : Quantization indexes  $q \mapsto$  reconstructed samples  $s'$

# Scalar Quantization: Discretization of Pdf



$$p_k = P(S' = s'_k) = \int_{u_k}^{u_{k+1}} f(s) ds$$

## Performance of Scalar Quantizers: Bit Rate



- Average bit rate  $R$  ( $\ell_k =$  codeword length for quantization index  $k$ )

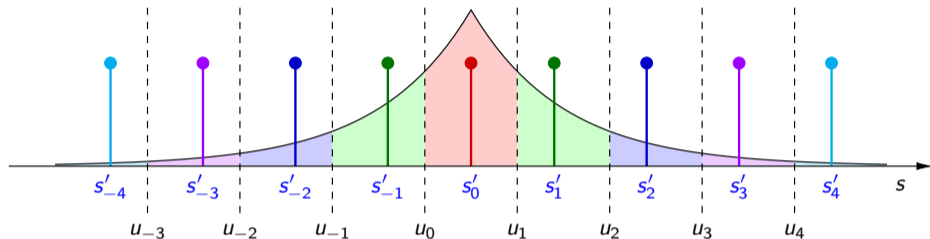
$$R = \mathbb{E}\{\ell(S')\} = \mathbb{E}\{\ell(\alpha(S))\} = \sum_k p_k \ell_k \quad \text{with} \quad p_k = \int_{u_k}^{u_{k+1}} f(s) ds$$

- Approximations (without knowledge of actual entropy coding)

→ fixed-length coding:  $R = \lceil \log_2 K \rceil$  ( $K$ : number of quantization intervals)

→ optimal entropy coding:  $R = H(S') = H(\alpha(S)) = - \sum_k p_k \log_2 p_k$

## Performance of Scalar Quantizers: MSE Distortion



- Average MSE distortion  $D$  is given by

$$D = \mathbb{E}\left\{ (S - Q(S))^2 \right\} = \int_{-\infty}^{\infty} (s - Q(s))^2 f(s) ds = \sum_{\forall k} \int_{u_k}^{u_{k+1}} (s - s'_k)^2 f(s) ds$$

- ➔ Similar for other additive distortion measures (e.g., all  $p$ -norm distortion measures)



# Optimal Scalar Quantizer for Fixed-Length Coding

## Goal: Minimize MSE Distortion for Quantizer with $K$ Quantization Intervals

- Neglect impact of entropy coding → Consider fixed-length coding

→ Rate  $R$  and MSE distortion  $D$  are given by

$$R = \lceil \log_2 K \rceil \quad (\text{typically } K = 2^B, \text{ with } B \text{ being the bits per codeword})$$

$$D = \sum_{\forall k} \int_{u_k}^{u_{k+1}} (s - s'_k)^2 f(s) ds$$

## Optimize Quantizer of size $K$

- Bit rate  $R$  is independent on decision thresholds and reconstruction levels ( $R$  is given by  $K$ )
- Distortion (MSE) depends on
  - $K$  reconstruction levels  $s'_k$
  - $K - 1$  decision thresholds  $u_k$

# Centroid Condition

$$D = \sum_{\forall i} \int_{u_i}^{u_{i+1}} (s - s'_i)^2 f(s) ds$$

- Optimize reconstruction levels  $s'_k$  for given decision thresholds  $u_k$

$$\frac{\partial}{\partial s'_k} D = \int_{u_k}^{u_{k+1}} 2 \cdot (s - s'_k) \cdot (-1) \cdot f(s) ds = 0$$

$$\int_{u_k}^{u_{k+1}} s f(s) ds = s'_k \cdot \int_{u_k}^{u_{k+1}} f(s) ds$$

## → Centroid Condition for MSE Distortion

$$s'_k = \mathbb{E}\{S \mid S \in \mathcal{I}_k\} = \frac{1}{p_k} \int_{u_k}^{u_{k+1}} s f(s) ds = \frac{\int_{u_k}^{u_{k+1}} s f(s) ds}{\int_{u_k}^{u_{k+1}} f(s) ds}$$

- Optimal reconstruction level  $s'_k$  is given by conditional mean

## Nearest Neighbour Condition

$$D = \sum_{\forall i} \int_{u_i}^{u_{i+1}} (s - s'_i)^2 f(s) ds$$

- Optimize decision thresholds  $u_k$  for given reconstruction levels  $s'_k$ 
  - Threshold  $u_k$  lies somewhere between neighboring reconstruction levels:  $s'_{k-1} < u_k < s'_k$
  - At the threshold  $u_k$ , we have the same distortion for both neighbouring intervals

$$(u_k - s'_{k-1})^2 = (u_k - s'_k)^2$$

$$u_k - s'_{k-1} = s'_k - u_k$$

$$2 u_k = s'_{k-1} + s'_k$$

### → Nearest Neighbour Condition for MSE Distortion

$$u_k = \frac{1}{2} (s'_{k-1} + s'_k)$$

- Optimal decision threshold  $u_k$  lies in the middle between the neighboring reconstruction levels

# Lloyd Quantizer: Minimization of Distortion

## Necessary Conditions for Minimizing MSE Distortion

1 Centroid condition

$$s'_k = \frac{\int_{u_k}^{u_{k+1}} s f(s) ds}{\int_{u_k}^{u_{k+1}} f(s) ds}$$

2 Nearest neighbour condition

$$u_k = \frac{1}{2}(s'_k + s'_{k-1})$$

## Design of Lloyd quantizers

- In general: Cannot be derived in closed form
- ➔ Iterative algorithm consisting of
  - Optimize decision thresholds  $u_k$  given reconstruction levels  $s'_k$
  - Optimize reconstruction levels  $s'_k$  given decision thresholds  $u_k$

## Lloyd Algorithm for Given Pdf (MSE Distortion)

- Given is:
- the size  $K$  of the quantizer (i.e., the number of quantization intervals)
  - the marginal probability density function  $f(s)$  of the source

### Iterative quantizer design

- 1 Choose an initial set of  $K$  reconstruction levels  $\{s'_k\}$
- 2 Update the  $K - 1$  decision thresholds  $\{u_k\}$  according to

$$u_k = \frac{s'_k + s'_{k-1}}{2} \quad (\text{nearest neighbor condition})$$

- 3 Update the  $K$  reconstruction levels  $\{s'_k\}$  according to

$$s'_k = \frac{\int_{u_k}^{u_{k+1}} s f(s) ds}{\int_{u_k}^{u_{k+1}} f(s) ds} \quad (\text{centroid condition})$$

- 4 Repeat the previous two steps until convergence

## Lloyd Algorithm for a Training Set (MSE Distortion)

- Given is:
- the size  $K$  of the quantizer (i.e., the number of quantization intervals)
  - a sufficiently large realization  $\{s_n\}$  of considered source

### Iterative quantizer design

- 1 Choose an initial set of  $K$  reconstruction levels  $\{s'_k\}$
- 2 Associate all samples of the training set  $\{s_n\}$  with one of the quantization intervals  $\mathcal{I}_k$

$$q(s_n) = \arg \min_{\forall k} (s_n - s'_k)^2 \quad (\text{nearest neighbor condition})$$

- 3 Update the reconstruction levels  $\{s'_k\}$  according to

$$s'_k = \frac{1}{N_k} \sum_{n: q(s_n)=k} s_n \quad (\text{centroid condition})$$

where  $N_k$  is the number of samples associated with  $\mathcal{I}_k$

- 4 Repeat the previous two steps until convergence

## Example: Lloyd Algorithm for Gaussian Source

### Gaussian Source

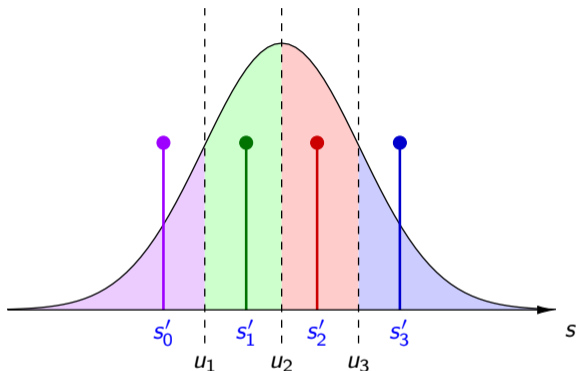
- Zero mean  $\mu = 0$
- Unit variance  $\sigma^2 = 1$

### Lloyd Quantizer of size $K = 4$

- Decision thresholds:
 

$u_1 = -0.982$
$u_2 = 0.000$
$u_3 = 0.982$
- Reconstruction levels:
 

$s'_0 = -1.510$
$s'_1 = -0.453$
$s'_2 = 0.453$
$s'_3 = 1.510$

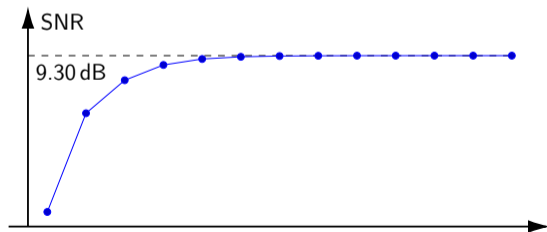
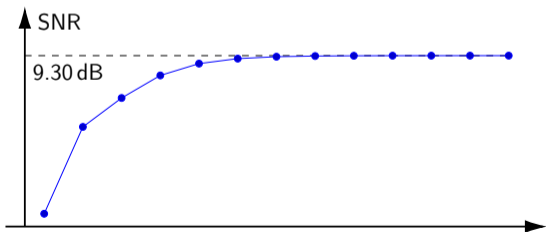
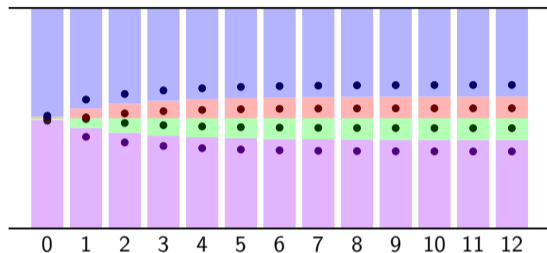
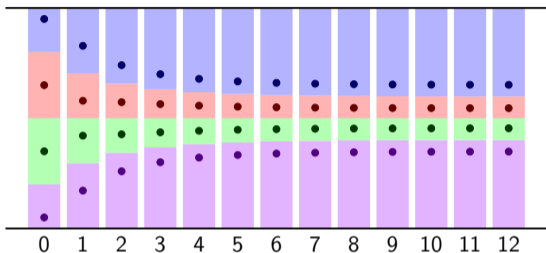


$$R = 2.0 \quad (\text{fixed-length coding})$$

$$D = 0.117$$

$$\text{SNR} = 9.30 \text{ dB}$$

# Example: Convergence of Lloyd Algorithm for Gaussian Source





## Example: Lloyd Algorithm for Laplacian Source

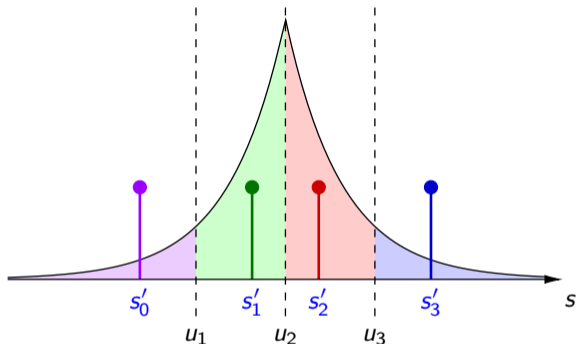
### Laplacian Source

- Zero mean  $\mu = 0$
- Unit variance  $\sigma^2 = 1$

### Lloyd Quantizer of size $K = 4$

- Decision thresholds:  $u_1 = -1.127$   
 $u_2 = 0.000$   
 $u_3 = 1.127$

- Reconstruction levels:  $s'_0 = -1.834$   
 $s'_1 = -0.420$   
 $s'_2 = 0.420$   
 $s'_3 = 1.834$

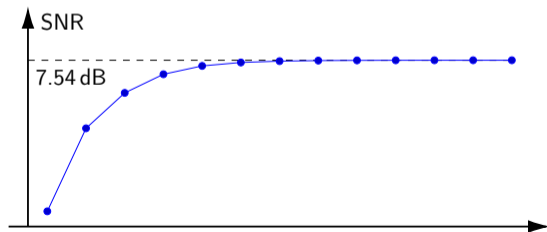
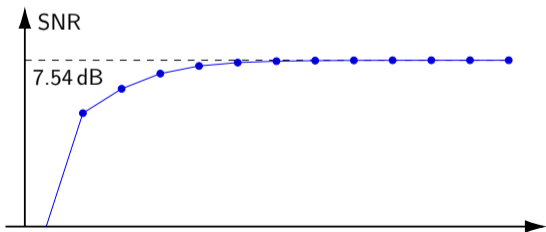
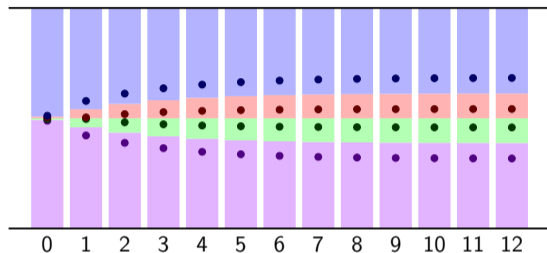
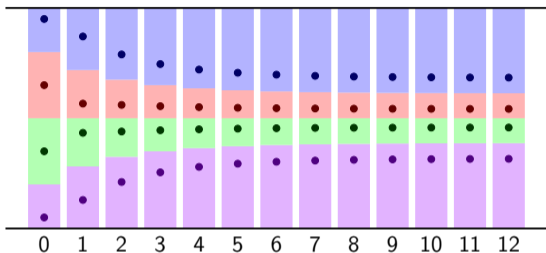


$$R = 2.0 \quad (\text{fixed-length coding})$$

$$D = 0.176$$

$$\text{SNR} = 7.54 \text{ dB}$$

# Example: Convergence of Lloyd Algorithm for Laplacian Source



# Centroid Quantizer at High Rates

## High-Rate Approximation

- High rates: Pdf  $f(s)$  is nearly constant inside each quantization interval

$$f(s) \approx \frac{p_k}{\Delta_k} = \frac{p_k}{u_{k+1} - u_k}$$

→ Direct consequence: Reconstruction value  $s'_k$  lies in center of quantization interval  $\mathcal{I}_k$

$$\begin{aligned} s'_k &= \frac{1}{p_k} \int_{u_k}^{u_{k+1}} s f(s) \, ds = \frac{1}{p_k} \cdot \frac{p_k}{u_{k+1} - u_k} \int_{u_k}^{u_{k+1}} s \, ds \\ &= \frac{1}{2} \cdot \frac{1}{u_{k+1} - u_k} \cdot (u_{k+1}^2 - u_k^2) = \frac{1}{2} \cdot \frac{(u_{k+1} + u_k) \cdot (u_{k+1} - u_k)}{u_{k+1} - u_k} \\ &= \frac{1}{2} (u_k + u_{k+1}) = u_k + \frac{\Delta_k}{2} \end{aligned}$$

# High-Rate Approximation of MSE Distortion for Centroid Quantizers

$$\begin{aligned}
 D &= \sum_{\forall k} \int_{u_k}^{u_{k+1}} (s - s'_k)^2 \cdot f(s) \, ds \\
 &= \sum_{\forall k} \frac{p_k}{\Delta_k} \int_{u_k}^{u_{k+1}} (s - s'_k)^2 \, ds \\
 &= \sum_{\forall k} \frac{p_k}{\Delta_k} \int_{u_k - s'_k}^{u_{k+1} - s'_k} t^2 \, dt \\
 &= \sum_{\forall k} \frac{p_k}{\Delta_k} \int_{-\Delta_k/2}^{\Delta_k/2} t^2 \, dt \\
 &= \sum_{\forall k} \frac{p_k}{\Delta_k} \cdot \frac{1}{3} \cdot \left( \frac{\Delta_k^3}{8} + \frac{\Delta_k^3}{8} \right)
 \end{aligned}$$

$$D = \frac{1}{12} \sum_{\forall k} p_k \Delta_k^2$$

## High-Rate Approximation: MSE Distortion for Lloyd Quantizer

- Will use: **Hölders inequality** in the following form (with  $x_k \geq 0$  and  $y_k \geq 0$ )

$$\alpha + \beta = 1 \quad \implies \quad \left( \sum_k x_k \right)^\alpha \cdot \left( \sum_k y_k \right)^\beta \geq \sum_k x_k^\alpha y_k^\beta$$

with equality iff  $y_k$  is proportional to  $x_k$ , i.e.,  $y_k = \text{const} \cdot x_k$

### Average MSE distortion of Lloyd quantizer of size $K$ (at high rates)

- Approximation for centroid quantizers

$$D = \frac{1}{12} \sum_{i=0}^{K-1} p_i \Delta_i^2 = \frac{1}{12} \sum_{i=0}^{K-1} f(s'_i) \Delta_i^3$$

- Rewrite expression using  $\sum_{i=0}^{K-1} (1/K) = K \cdot (1/K) = 1$

$$D = \frac{1}{12} \left( \left( \sum_{i=0}^{K-1} f(s'_i) \Delta_i^3 \right)^{\frac{1}{3}} \cdot \left( \sum_{i=0}^{K-1} \frac{1}{K} \right)^{\frac{2}{3}} \right)^3$$

## High-Rate Approximation: MSE Distortion for Lloyd Quantizer

**Average MSE distortion of Lloyd quantizer of size  $K$**  (at high rates)

- Apply Hölders inequality

$$D \geq \frac{1}{12} \left( \sum_{i=0}^{K-1} \left( f(s'_i) \Delta_i^3 \right)^{\frac{1}{3}} \left( \frac{1}{K} \right)^{\frac{2}{3}} \right)^3 = \frac{1}{12 K^2} \left( \sum_{i=0}^{K-1} \sqrt[3]{f(s'_i)} \Delta_i \right)^3$$

with equality iff  $\Delta_i \sqrt[3]{f(s'_i)} = \text{const}$

- Remember: Lloyd quantizer minimizes distortion for given size  $K$

$$D = \frac{1}{12 K^2} \left( \sum_{i=0}^{K-1} \sqrt[3]{f(s'_i)} \Delta_i \right)^3$$

- Asymptotic limit for large  $K$  ( $\Delta_k \rightarrow 0$ )

$$D = \frac{1}{12 K^2} \left( \int_{-\infty}^{\infty} \sqrt[3]{f(s)} ds \right)^3$$

## High-Rate Approximation: Lloyd Quantizer with Fixed-Length Coding

### MSE Distortion for Lloyd Quantizer at High Rates and Rate for Fixed-Length Coding

$$D = \frac{1}{12 K^2} \left( \int_{-\infty}^{\infty} \sqrt[3]{f(s)} \, ds \right)^3 \quad \text{and} \quad R = \log_2 K \quad \Longrightarrow \quad \frac{1}{K^2} = 2^{-2R}$$

### Lloyd Quantizer with Fixed-Length Coding at High Rates

- Panter and Dite approximation for operational distortion-rate function

$$D_F(R) = \frac{1}{12} \left( \int_{-\infty}^{\infty} \sqrt[3]{f(s)} \, ds \right)^3 \cdot 2^{-2R}$$

$$D_F(R) = \varepsilon_F^2 \cdot \sigma^2 \cdot 2^{-2R} \quad \text{with}$$

$$\varepsilon_F^2 = \frac{1}{12 \sigma^2} \left( \int_{-\infty}^{\infty} \sqrt[3]{f(s)} \, ds \right)^3$$

## Lloyd Quantizer with Fixed-Length Coding vs Panter-Dite Approximation

Uniform pdf

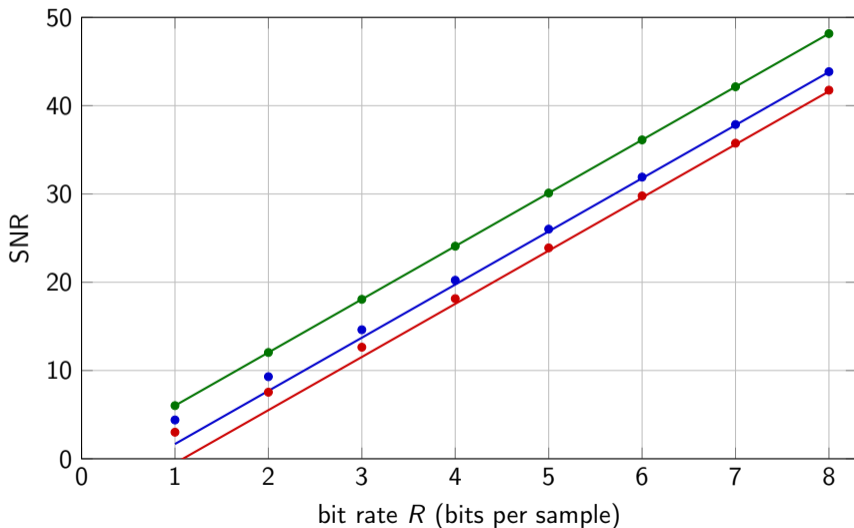
$$\varepsilon_F^2 = 1$$

Gaussian pdf

$$\varepsilon_F^2 = \frac{\sqrt{3}\pi}{2}$$

Laplacian pdf

$$\varepsilon_F^2 = 4.5$$





# Summary of Lecture

## Scalar Quantization

- Input-output function  $s' = Q(s)$  is a staircase function
- Quantizer is characterized by  $K$  reconstruction levels  $s'_k$  and  $K - 1$  decision thresholds  $u_k$

## Lloyd Quantizer

- Minimizes distortion  $D$  for given number  $K$  of quantization intervals
- Two optimization criterions
  - Centroid condition (MSE):  $s'_k = E\{S \mid S \in \mathcal{I}_k\}$
  - Nearest neighbor condition (MSE):  $u_k = (s'_k + s'_{k-1})/2$
- Lloyd quantizer design: Iterate between the two optimization criterions
- High-rate approximation of Lloyd quantizer with fixed-length coding (Panter-Dite approximation)

## Next Steps

- Theoretical limits for lossy source coding
- Consider entropy coding in quantizer design

## Exercise 1: Implement Lloyd Algorithm

Implement the Lloyd algorithm using a programming language of your choice.

- Test the algorithm (for quantizer sizes of  $K = 2, 4, 8, 16, 32$ ) for

- a unit-variance Gaussian pdf:

$$f(s) = \frac{1}{\sqrt{2\pi}} e^{-\frac{1}{2} s^2}$$

- a unit-variance Laplacian pdf:

$$f(s) = \frac{1}{\sqrt{2}} e^{-\sqrt{2} |s|}$$

- Determine the distortion  $D$  for your quantizers.
- Compare the R-D performance of your quantizers (for  $K = 2, 4, 8, 16, 32$ ) to the high-rate approximation for Lloyd quantizers with fixed-length codes.

You can implement the Lloyd algorithm that directly uses the pdf or the Lloyd algorithm that uses a training set (files with 1 000 000 samples in float32 format are provided on the course web site)

## Exercise 2: Lloyd Quantizer for MSE Distortion (Alternative)

Given is a stationary source with a zero-mean Laplace pdf  $f(x)$  and a symmetric 3-interval quantizer:

$$f(x) = \frac{1}{\sqrt{2\sigma^2}} e^{-\sqrt{\frac{2}{\sigma^2}}|x|} \quad \text{and} \quad Q(x) = \begin{cases} -b & : x < -a \\ 0 & : |x| \leq a \\ b & : x > a \end{cases}$$

- (a) Derive the optimal reconstruction value  $b$  as a function of the threshold  $a$  for MSE distortion. Express the resulting distortion as function of the threshold  $a$  and the variance  $\sigma^2$ .
- (b) Determine the decision threshold  $a$  in a way that a Lloyd quantizer for MSE distortion is obtained. Determine the distortion and rate for the Lloyd quantizer by assuming fixed-length coding ( $R = \log_2 K$ ) and compare the obtained R-D point with the high-rate approximation.
- (c) Can the derived optimal quantizer for fixed-length coding be improved by adding entropy coding (without changing the decision thresholds and reconstruction levels)?

## Exercise 3: Lloyd Quantizer for MAE Distortion (Another Alternative)

Given is a stationary source with a zero-mean Laplace pdf  $f(x)$  and a symmetric 3-interval quantizer:

$$f(x) = \frac{1}{2m} e^{-\frac{|x|}{m}} \quad \text{and} \quad Q(x) = \begin{cases} -b & : & x < -a \\ 0 & : & |x| \leq a \\ b & : & x > a \end{cases}$$

- (a) Derive the centroid condition and nearest neighbor condition for MAE distortion

$$D = E\{|S - S'|\}$$

- (b) Derive the optimal reconstruction value  $b$  as a function of the threshold  $a$  for MAE distortion. Express the resulting distortion as function of the threshold  $a$  and the parameter  $m$ .
- (c) Determine the decision threshold  $a$  in a way that a Lloyd quantizer for MAE distortion is obtained. Determine the distortion and rate for the quantizer by assuming fixed-length coding ( $R = \log_2 K$ ).