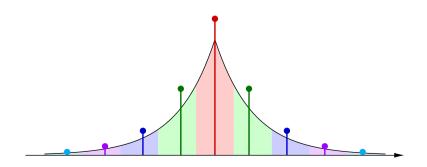
Scalar Quantization



Last Lectures: Lossless Coding

Variable-length Coding

- Scalar codes, conditional codes, block codes, V2V codes (using codeword tables)
- For given pmf: **Huffman algorithm** yields optimal codeword table
- Problem: Codeword tables become too large for practical application of block codes

Arithmetic Coding

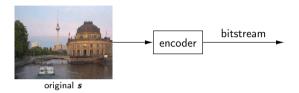
- No codeword table: On-the-fly encoding and decoding
- lacksquare Sub-optimal block code for arbitrarily large block sizes N (very close to optimum for $N\gg 1$)
- Straightforward combination with conditional and adaptive probability models

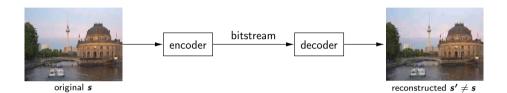
Reduction of Inter-Symbols Dependencies before Entropy Coding

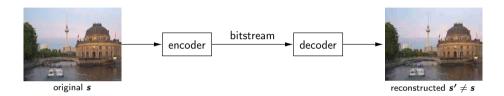
- Affine and linear prediction: Suitable for reducing dependencies in audio, image, video data
- Lempel-Ziv coding or block sorting: Suitable for text, source code, general files
- → Lossless coding in practice: Prediction followed by entropy coding of prediction errors
 - Lempel-Ziv coding or block sorting followed by entropy coding



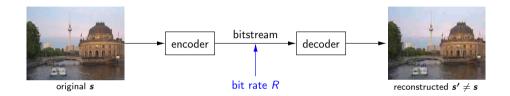
original $m{s}$





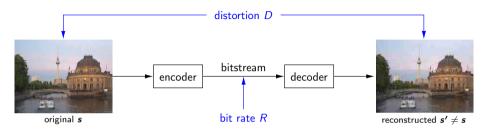


Lossy coding is characterized by two aspects:



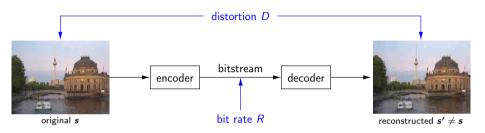
Lossy coding is characterized by two aspects:

■ Bit rate R: Average number of bits per sample (or per time unit)



Lossy coding is characterized by two aspects:

- Bit rate R: Average number of bits per sample (or per time unit)
- Distortion D: Measure for deviation between original signal s and reconstructed signal s'

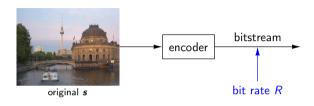


Lossy coding is characterized by two aspects:

- Bit rate R: Average number of bits per sample (or per time unit)
- Distortion D: Measure for deviation between original signal s and reconstructed signal s'

Design Goal: Smallest possible bit rate for given maximum distortion, or Smallest possible distortion for given maximum bit rate

Lossy Coding: Bit Rate



Bit Rate R:

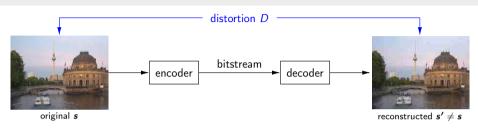
■ Images: Average number of bits per sample

Audio or video: Average number of bits per time units

Often used Approximation:

- Assume that we have a close to optimal entropy coding (e.g., arithmetic coding)
- Bit rate = Entropy of symbols that are actually transmitted

Lossy Coding: Distortion



Distortion Measures used in Practice

■ General *p*-norm distortion:

$$D_p = rac{1}{N} \sum_{k=1}^{N} \left| s_k - s_k' \right|^p$$
 or $D_p = \mathrm{E} \left\{ \left| S - S' \right|^p \right\}$

■ Most often: Mean squared error (MSE)

$$D_2 = \frac{1}{N} \sum_{k=1}^{N} (s_k - s'_k)^2$$
 or $D_2 = \mathbb{E}\left\{ (S - S')^2 \right\}$

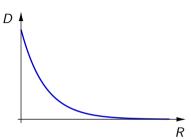
MSE Distortion as Signal-To-Noise Ratio (SNR)

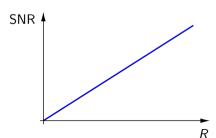
Signal-to-Noise Ratio (SNR)

■ Logarithmic ratio of variance and MSE distortion

$$\mathsf{SNR} = 10 \cdot \mathsf{log}_{10} igg(rac{\sigma^2}{D_2} igg)$$

Measured in decibel (dB)





MSE Distortion as Signal-To-Noise Ratio (SNR)

Signal-to-Noise Ratio (SNR)

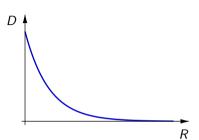
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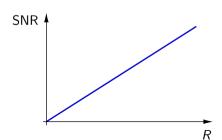
$$\mathsf{SNR} = 10 \cdot \mathsf{log_{10}} igg(rac{\sigma^2}{D_2} igg)$$

Measured in decibel (dB)

Advantages of using SNR

Independent of signal variance





MSE Distortion as Signal-To-Noise Ratio (SNR)

Signal-to-Noise Ratio (SNR)

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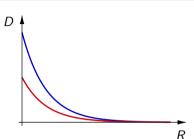
Measured in decibel (dB)

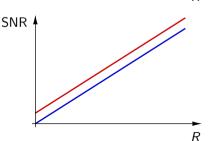
Advantages of using SNR

- Independent of signal variance
- Easy interpretation of differences

$$\Delta \mathsf{SNR} = \mathsf{SNR}_a - \mathsf{SNR}_b = -10 \cdot \log_{10} rac{D_a}{D_b}$$

■ Examples: $D_a = D_b$ \rightarrow $\Delta SNR \approx 0.0 \, dB$ $D_a = D_b/\sqrt{2}$ \rightarrow $\Delta SNR \approx 1.5 \, dB$ $D_a = D_b/2$ \rightarrow $\Delta SNR \approx 3.0 \, dB$ $D_a = D_b/4$ \rightarrow $\Delta SNR \approx 6.0 \, dB$ $D_a = D_b/8$ \rightarrow $\Delta SNR \approx 9.0 \, dB$





Probabilistic Modeling of Sources

Source Coding in Practice

- Encoder and decoder are computer programs
- → Actual input signals are **discrete-time** and **discrete-amplitude** signals

Real-world signals

- In most cases: Continuous-time and continuous-amplitude signals
- Discrete signals are obtained by sampling and quantization
- Typical scenarios: Initial quantization has negligible effect on source coding

Probabilistic Modeling of Sources

Source Coding in Practice

- Encoder and decoder are computer programs
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Real-world signals

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- Discrete signals are obtained by sampling and quantization
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Theoretical Analysis of Lossy Source Coding

- Will mostly use models for **discrete-time** and **continuous-amplitude** signals
- Main reason: Mathematical tractability
- → Interpretation: Consider signal before initial quantization

Review: Random Variables and Cumulative Distribution Function (CDF)

Random Variable

■ Function $X(\zeta)$ of the sample space \mathcal{O} that assigns a real value $x = X(\zeta)$ to each possible outcome $\zeta \in \mathcal{O}$ of a random experiment

Cumulative Distribution Function (cdf)

Cumulative distribution function $F_X(x)$ of a random variable X

$$F_X(s) = P(X \le x) = P(\lbrace \zeta : X(\zeta) \le x \rbrace)$$

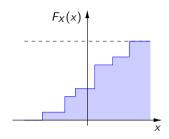
■ Joint cdf of two random variables X and Y

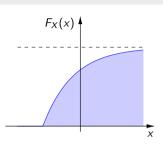
$$F_{XY}(x,y) = P(X \le x, Y \le y)$$

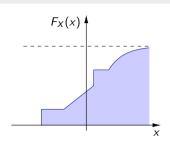
■ Conditional cdf of a random variable X given another random variable Y

$$F_{X|Y}(x|y) = P(X \le x \mid Y \le y) = \frac{P(X \le x, Y \le y)}{P(Y \le y)} = \frac{F_{XY}(x, y)}{F_{Y}(y)}$$

Review: Examples of Cumulative Distribution Functions







Staircase function

- Random variable X can only take a countable number of values
- → Discrete random variable

Continuous function

- Random variable X can take all values inside one or more non-zero intervals
- → Continuous random variable

Mixed type

■ Random variable X can take all values inside one or more non-zero intervals and a countable number of additional values

Continuous Random Variables and Probability Density Function (PDF)

Continuous Random Variables

■ A random variable X is called a **continuous random variable** if and only if its cdf $F_X(x)$ is a continuous function

Probability Density Function

 \blacksquare Probability density function (pdf) of a continuous random variable S

$$f_X(x) = \frac{\partial}{\partial x} F_X(x) \quad \Longleftrightarrow \quad F_X(x) = \int_{-\infty}^{x} f_X(t) dt$$

- → Properties: $f_X(x) \ge 0, \forall x$
 - $\int_{-\infty}^{\infty} f_X(t) dt = 1$
 - $P(a < X \le b) = \int_a^b f_X(t) dt$

Examples for Continuous Distributions (Zero Mean)

Uniform

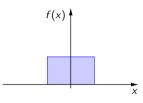
$$f(x) = \begin{cases} \frac{1}{2a} : |x| \le a \\ 0 : \text{ otherwise} \end{cases}$$

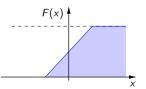
Gaussian

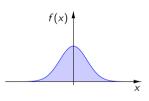
$$f(x) = \frac{1}{\sqrt{2\pi\sigma^2}} e^{-\frac{(x-\mu)^2}{2\sigma^2}}$$

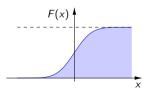
Laplacian

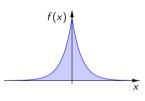
$$f(x) = \frac{1}{\sqrt{2\sigma^2}} e^{-\sqrt{\frac{2}{\sigma^2}}|x-\mu|}$$

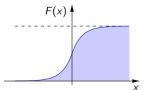






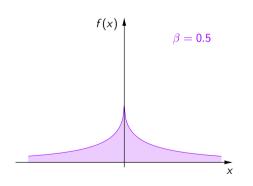


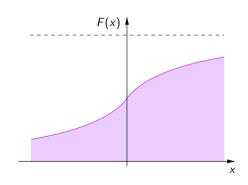




$$f(x) = \frac{\beta}{2 \alpha \Gamma(1/\beta)} e^{-\left(\frac{|x-\mu|}{\alpha}\right)^{\beta}} \quad \text{with} \quad \Gamma(x) = \int_0^{\infty} e^{-t} t^{x-1} dt$$

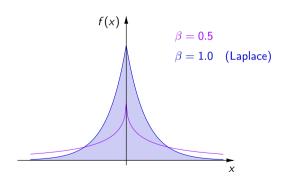
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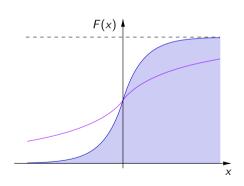




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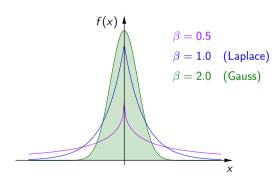
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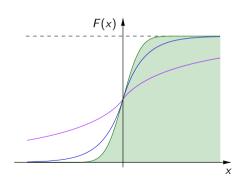




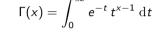
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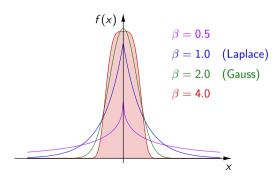
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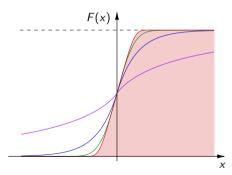




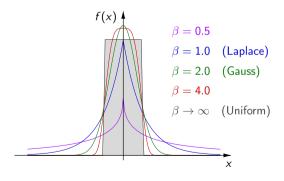
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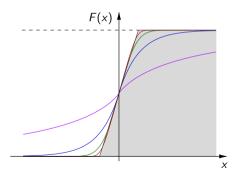






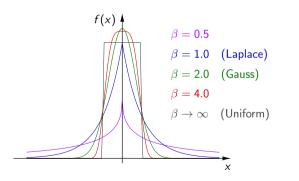
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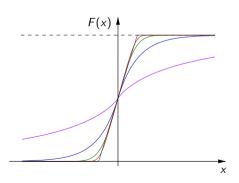




Shape parameter $\beta \in (0, \infty)$:

$$f(x) = \frac{\beta}{2 \alpha \Gamma(1/\beta)} e^{-\left(\frac{|x-\mu|}{\alpha}\right)^{\beta}} \quad \text{with} \quad \Gamma(x) = \int_0^{\infty} e^{-t} t^{x-1} dt$$





Suitable approximation for many distributions

Joint and Conditional Probability Density Function

Joint Probability Density Function

■ Joint pdf of two random variables X and Y

$$f_{XY}(x,y) = \frac{\partial^2}{\partial x \, \partial y} \, F_{XY}(x,y)$$

Joint and Conditional Probability Density Function

Joint Probability Density Function

■ Joint pdf of two random variables X and Y

$$f_{XY}(x,y) = \frac{\partial^2}{\partial x \, \partial y} \, F_{XY}(x,y)$$

Conditional Probability Density Function

Conditional pdf of a random variable X given another random variable Y

$$f_{X|Y}(x|y) = \frac{\partial}{\partial x} F_{X|Y}(x|y) = \frac{\partial}{\partial x} \frac{F_{XY}(x,y)}{F_{Y}(y)} = \frac{\frac{\partial^{2}}{\partial x \partial y} F_{XY}(x,y)}{\frac{\partial}{\partial y} F_{Y}(y)} = \frac{f_{XY}(x,y)}{f_{Y}(y)}$$

Expected Values for Continuous Random Variables

Expected Values

Expected value of a function g(X) of a continuous random variable X

$$\mathrm{E}\{g(X)\} = \mathrm{E}_X\{g(X)\} = \int_{-\infty}^{\infty} g(x) f_X(x) \, \mathrm{d}x$$

Expected value of function g(X, Y) of two continuous random variables X and Y

$$\mathrm{E}\{g(X,Y)\} = \mathrm{E}_{XY}\{g(X,Y)\} = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} g(x,y) f_{XY}(x,y) \, \mathrm{d}x \, \mathrm{d}y$$

Conditional Expected Values

Expected value of function g(X) of a random variable X given another random variable Y

$$E\{g(X) \mid Y\} = \int_{-\infty}^{\infty} g(x) f_{X|Y}(x \mid Y) dx$$
 (is another random variable)

Properties of Expected Values

same properties as in discrete case

Important Properties

■ Linearity of expected values

$$E\{aX + bY\} = a \cdot E\{X\} + b \cdot E\{Y\}$$

 \blacksquare For independent random variables X and Y

$$E\{XY\} = E\{X\} E\{Y\}$$

Iterative expectation rule

$$E\{ E\{ g(X) | Y \} \} = E\{ g(X) \}$$

Important Expected Values

■ Mean μ_X of a random variable X

$$\mu_X = \mathrm{E}\{X\} = \int_{-\infty}^{\infty} x \cdot f_X(x) \, \mathrm{d}x$$

■ Variance σ_X^2 of a random variable X

$$\sigma_X^2 = \mathrm{E}\Big\{ (X - \mathrm{E}\{X\})^2 \Big\} = \int_{-\infty}^{\infty} (x - \mu_X)^2 \cdot f_X(x) \, \mathrm{d}x$$

Covariance σ_{XY}^2 of two random variables X and Y, and correlation coefficient ϕ_{XY}

$$\sigma_{XY}^{2} = \mathrm{E}\left\{ (X - \mathrm{E}\{X\}) (Y - \mathrm{E}\{Y\}) \right\} = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} (x - \mu_{x})(y - \mu_{y}) \cdot f_{XY}(x, y) \, \mathrm{d}x \, \mathrm{d}y$$

$$\phi_{XY} = \frac{\sigma_{XY}^{2}}{\sqrt{\sigma_{X}^{2} \cdot \sigma_{Y}^{2}}}$$

Continuous Random Processes

Discrete-Time Random Process

- Series of random experiments at time instants t_n , with $n = 0, 1, 2, \cdots$
- For each experiment: Random variable $X_n = X(t_n)$
- Random process: Series of random variables

$$X = \{X_0, X_1, X_2, \cdots\} = \{X_n\}$$

Discrete-Time Continuous-Amplitude Random Process

- \blacksquare Random variables X_n are continuous random variables
- → Type of random processes we consider for analyzing lossy coding

Statistical Properties of Continuous Random Processes

Characterization of Statistical Properties

■ Consider *N*-dimensional random vector

$$\boldsymbol{X}_{k}^{(N)} = \{X_{k}, X_{k+1}, \cdots, X_{k+N-1}\}$$

■ N-th order joint cdf

$$F_k^{(N)}(\mathbf{x}) = P(\mathbf{X}_k^{(N)} \le \mathbf{x}) = P(X_k \le x_0, X_{k+1} \le x_1, \cdots, X_{k+N-1} \le x_{N-1})$$

N-th order joint pdf

$$f_k^{(N)}(\mathbf{x}) = \frac{\partial^N}{\partial x_0 \cdots \partial x_{N-1}} F_k^{(N)}(\mathbf{x})$$

■ Also: Conditional cdfs and conditional pdfs

Models for Random Processes

Stationary Random Processes

- Statistical properties are invariant to a shift in time
- In this course: Typically restrict our considerations to stationary processes

Memoryless Random Processes

 \blacksquare All random variables X_n are independent of each other

Independent and Identically Distributed (IID) Random Processes

■ Random processes that are stationary and memoryless

Markov Processes

■ Markov property: Future outcomes do only depend on present outcome, but not on past outcomes

$$F(x_n | x_{n-1}, x_{n-2}, x_{n-3}, \cdots) = F(x_n | x_{n-1})$$

$$f(x_n | x_{n-1}, x_{n-2}, x_{n-3}, \cdots) = f(x_n | x_{n-1})$$

Simple model for random processes with memory

Autoregressive (AR) Processes

General AR(p) Model

■ Autoregressive model of order p for random variables X_n with mean μ

$$X_n = Z_n + \mu + \sum_{k=1}^p \varrho_k \cdot (X_{n-k} - \mu)$$

where $\mathbf{Z} = \{Z_n\}$ is a zero-mean iid process (innovation process) and $\varrho_1, \dots, \varrho_p$ are the model parameters

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Special case: AR(1) model

■ Autoregressive model of order p = 1

$$X_n = Z_n + \mu + \varrho \cdot (X_{n-1} - \mu)$$

Autoregressive (AR) Processes

General AR(p) Model

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ightharpoonup Completely specified by mean μ , correlation coefficient ϱ , and pdf $f_Z(z)$ of iid process $\{Z_n\}$

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Special case: AR(1) model

■ Autoregressive model of order p = 1

$$X_n = Z_n + \mu + \varrho \cdot (X_{n-1} - \mu)$$

- \rightarrow Completely specified by mean μ , correlation coefficient ϱ , and pdf $f_Z(z)$ of iid process $\{Z_n\}$
- → Important type of stationary Markov process for continuous random processes

Gaussian Processes

Gaussian Random Process

- \blacksquare All finite collections of random variables X_n are Gaussian random vectors
- N-th order pdf is given by N-th order auto-covariance matrix C_N and mean μ

$$f_{\mathbf{X}}(\mathbf{x}) = \frac{1}{\sqrt{(2\pi)^N |\mathbf{C}_N|}} e^{-\frac{1}{2}(\mathbf{x}-\boldsymbol{\mu})^{\mathrm{T}} \mathbf{C}_N^{-1}(\mathbf{x}-\boldsymbol{\mu})} \quad \text{with} \quad \boldsymbol{\mu} = \begin{pmatrix} \mu \\ \vdots \\ \mu \end{pmatrix}$$

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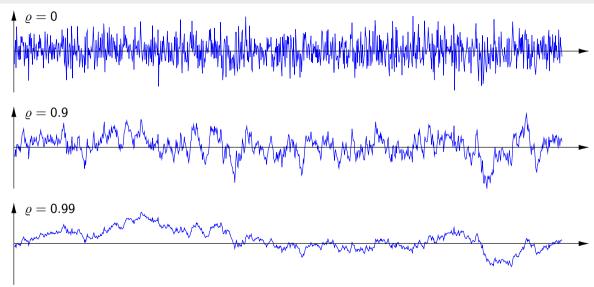
Stationary Gauss-Markov Process

- Stationary Markov process that is also a Gaussian random process
- Can be constructed with Gaussian iid process $Z = \{Z_n\}$ according to

$$X_n = \mu + \varrho \left(X_{n-1} - \mu \right) + Z_n$$

- Statistical properties are completely specified by mean μ , variance σ^2 , correlation coefficient ϱ
- → Will use it as very simple model for sources with memory

Examples of Gauss-Markov Processes (1000 Samples)



Summary of Mathematical Basics (for Continuous Case)

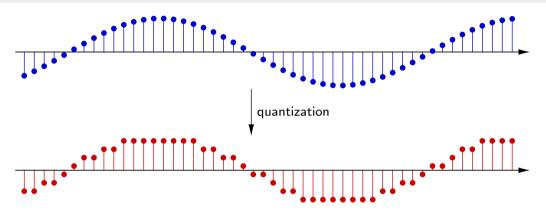
Continuous Random Variables

- Can take all values inside one or more non-zero intervals
- Cumulative distribution function (cdf): Continuous function
- Probability density function (pdf)
- Expected values: Mean, variance, covariance

Discrete-Time Continuous-Amplitude Random Processes

- Sequence of continuous random variables: Model for lossy source coding
- Types of random processes: Stationary, memoryless, iid, Markov
- Suitable model for real signals: Autoregressive processes
- Special importance for lossy source coding: Gaussian processes
- Simple model for sources with memory: Gauss-Markov process

Quantization



- "Lossy part" of source coding
- Non-reversible mapping from input to output samples
- Determines trade-off between signal fidelity and bit rate

Scalar Quantization: Functional Mapping

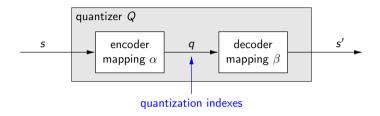


Scalar Quantization: Functional mapping of an input sample to an output sample

$$s'=Q(s)$$

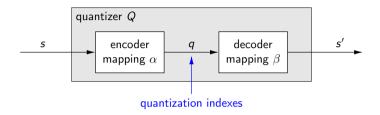
- Input: Discrete or continuous
- Output: Set of obtainable output points is countable
 - Less obtainable output points than input points
- → Non-reversible loss in signal fidelity

Structure of Scalar Quantizers: Encoder and Decoder Mapping



lacksquare Split quantizer Q into encoder mapping lpha and decoder mapping eta

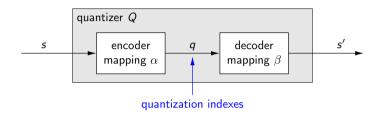
Structure of Scalar Quantizers: Encoder and Decoder Mapping



- Split quantizer Q into encoder mapping α and decoder mapping β
- **E**ncoder mapping α : Maps input sample s to a quantizer index q (integer)

$$q = \alpha(s)$$

Structure of Scalar Quantizers: Encoder and Decoder Mapping

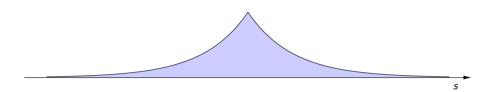


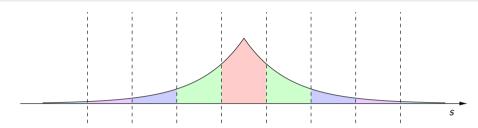
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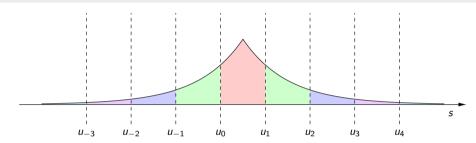
■ Decoder mapping β : Maps quantizer index q to reconstructed samples s'

$$s' = \beta(q) = \beta(\alpha(s)) = Q(s)$$

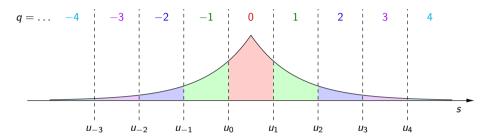




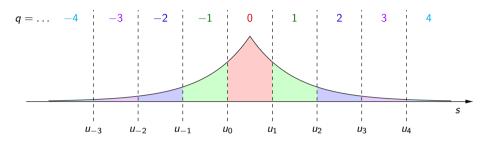
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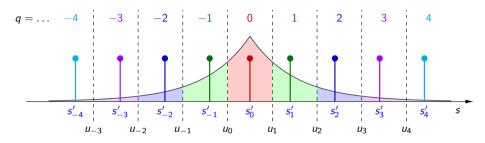
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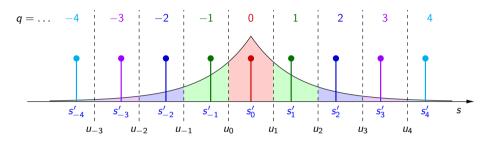
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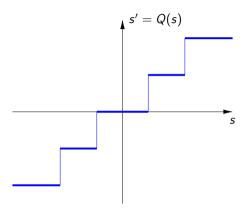
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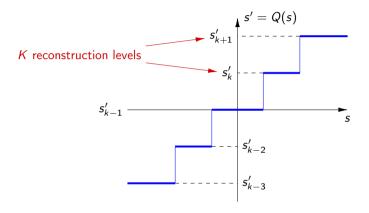


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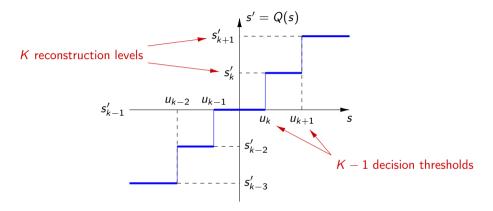


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- \rightarrow Scalar quantization: Replace input value s that falls inside \mathcal{I}_k with reconstruction value s'_k

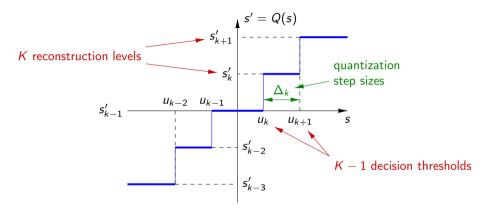




■ Scalar quantizer mapping:
$$Q: \mathbb{R} \mapsto \{\cdots, s'_{k-1}, s'_k, s'_{k+1}, \cdots\}$$



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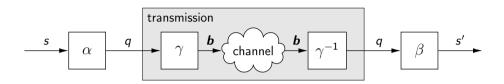
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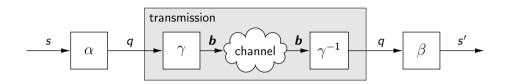
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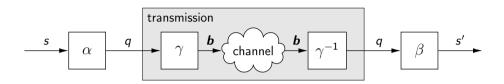
- Quantization step sizes:
- $\Delta_k = u_{k+1} u_k$



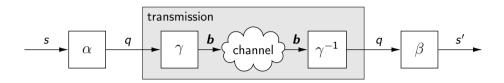
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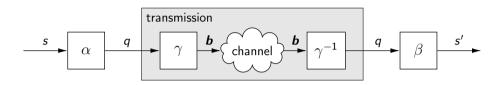
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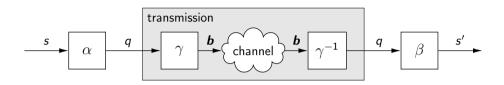
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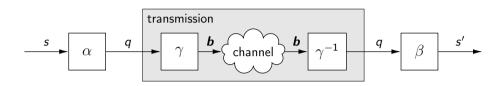
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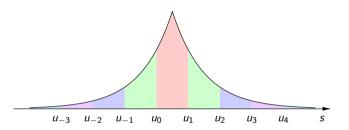


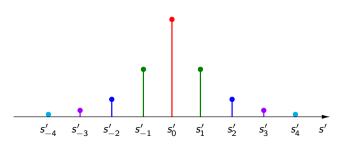
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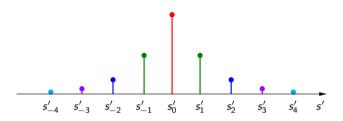
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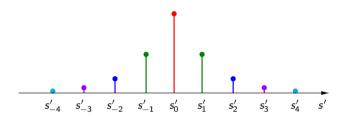
Scalar Quantization: Discretization of Pdf





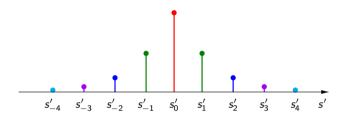
$$p_k = P(S' = s'_k) = \int_{u_k}^{u_{k+1}} f(s) ds$$





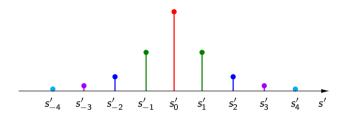
■ Average bit rate R (ℓ_k = codeword length for quantization index k)

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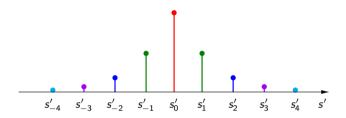
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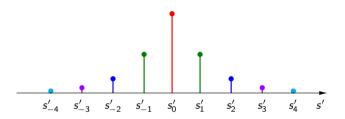
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- Approximations (without knowledge of actual entropy coding)

$$R = \lceil \log_2 K \rceil$$

 \rightarrow fixed-length coding: $R = \lceil \log_2 K \rceil$ (K: number of quantization intervals)

Performance of Scalar Quantizers: Bit Rate



• Average bit rate R (ℓ_k = codeword length for quantization index k)

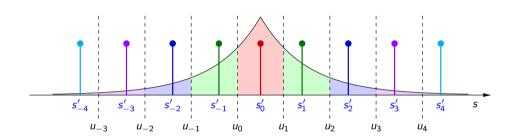
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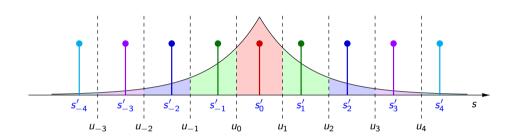
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$$ightharpoonup$$
 optimal entropy coding: $R = H(S') = H(\alpha(S)) = -\sum p_k \log_2 p_k$



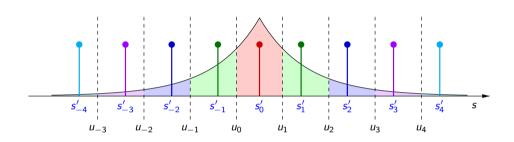
■ Average MSE distortion *D* is given by

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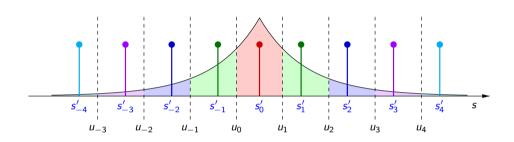
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→ Similar for other additive distortion measures (e.g., all p-norm distortion measures)

Optimal Scalar Quantizer for Fixed-Length Coding

Goal: Minimize MSE Distortion for Quantizer with K Quantization Intervals

- Neglect impact of entropy coding
 → Consider fixed-length coding
- \rightarrow Rate R and MSE distortion D are given by

$$R=\lceil \log_2 K
ceil$$
 (typically $K=2^B$, with B being the bits per codeword) $D=\sum_{orall k}\int_{u_k}^{u_{k+1}}(s-s_k')^2\,f(s)\,\mathrm{d} s$

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- Distortion (MSE) depends on
 - \rightarrow K reconstruction levels s'_k
 - $\rightarrow K-1$ decision thresholds u_k

$$D = \sum_{\forall i} \int_{u_i}^{u_{i+1}} (s - s_i')^2 f(s) ds$$

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→ Centroid Condition for MSE Distortion

$$\boxed{s'_k = \mathrm{E}\{\,S \,|\, S \in \mathcal{I}_k\,\} = \frac{1}{p_k} \int_{u_k}^{u_{k+1}} s \,f(s) \,\mathrm{d}s = \frac{\int_{u_k}^{u_{k+1}} s \,f(s) \,\mathrm{d}s}{\int_{u_k}^{u_{k+1}} f(s) \,\mathrm{d}s}}$$

 \rightarrow Optimal reconstruction level s'_{ν} is given by conditional mean

$$D = \sum_{\forall j} \int_{u_j}^{u_{j+1}} (s - s_j')^2 f(s) ds$$

lacksquare Optimize decision thresholds u_k for given reconstruction levels s_k'

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→ Nearest Neighbour Condition for MSE Distortion

$$u_k = \frac{1}{2} \left(s'_{k-1} + s'_k \right)$$

 \rightarrow Optimal decision threshold u_k lies in the middle between the neighboring reconstruction levels

Necessary Conditions for Minimizing MSE Distortion

1 Centroid condition

$$s_k' = \frac{\int_{u_k}^{u_{k+1}} s f(s) ds}{\int_{u_k}^{u_{k+1}} f(s) ds}$$

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Design of Lloyd quantizers

- In general: Cannot be derived in closed form
- → Iterative algorithm consisting of
 - Optimize decision thresholds u_k given reconstruction levels s'_k
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 (centroid condition)

4 Repeat the previous two steps until convergence

Given is:

- the size K of the quantizer (i.e., the number of quantization intervals)
- a sufficiently large realization $\{s_n\}$ of considered source

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Iterative quantizer design

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Iterative quantizer design

- 1 Choose an initial set of K reconstruction levels $\{s'_k\}$
- **2** Associate all samples of the training set $\{s_n\}$ with one of the quantization intervals \mathcal{I}_k

$$q(s_n) = \arg\min_{\forall k} (s_n - s'_k)^2$$
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4 Repeat the previous two steps until convergence

Example: Lloyd Algorithm for Gaussian Source

Gaussian Source

- \blacksquare Zero mean $\mu=0$
- Unit variance $\sigma^2 = 1$

Lloyd Quantizer of size K = 4

■ Decision thresholds: $u_1 = -0.982$

$$u_1 = -0.962$$
 $u_2 = 0.000$

$$u_2 = 0.000$$

$$u_3 = 0.982$$

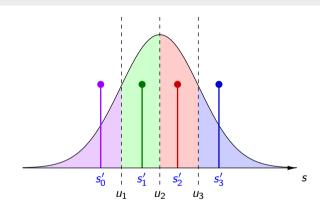
■ Reconstruction levels:

$$s_0' = -1.510$$

$$s_1' = -0.453$$

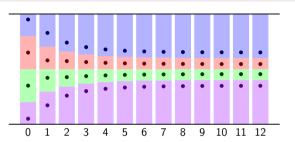
$$s_2' = 0.453$$

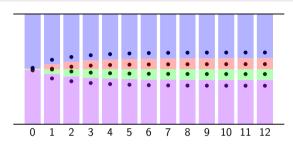
$$s_3' = 1.510$$

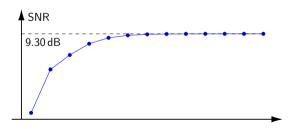


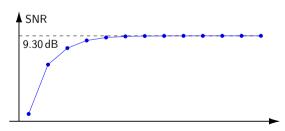
$$R = 2.0$$
 (fixed-length coding)
 $D = 0.117$
SNR = 9.30 dB

Example: Convergence of Lloyd Algorithm for Gaussian Source









Example: Lloyd Algorithm for Laplacian Source

Laplacian Source

- \blacksquare Zero mean $\mu=0$
- Unit variance $\sigma^2 = 1$

Lloyd Quantizer of size K = 4

■ Decision thresholds: $u_1 = -1.127$

$$u_1 = -1.127$$
 $u_2 = 0.000$

$$u_2 = 0.000$$

$$u_3 = 1.127$$

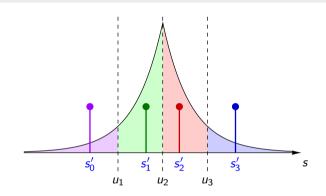
■ Reconstruction levels:

$$s_0' = -1.834$$

$$s_1' = -0.420$$

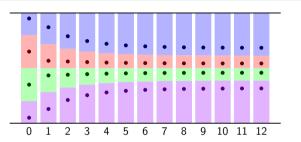
$$s_2' = 0.420$$

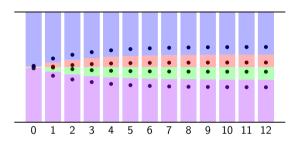
$$s_3' = 1.834$$

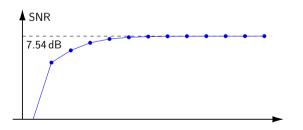


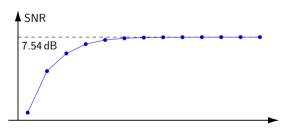
$$R = 2.0$$
 (fixed-length coding)
 $D = 0.176$
SNR = 7.54 dB

Example: Convergence of Lloyd Algorithm for Laplacian Source









Centroid Quantizer at High Rates

High-Rate Approximation

■ High rates: Pdf f(s) is nearly constant inside each quantization interval

$$f(s) pprox rac{p_k}{\Delta_k} = rac{p_k}{u_{k+1} - u_k}$$

 \rightarrow Direct consequence: Reconstruction value s'_{k} lies in center of quantization interval \mathcal{I}_{k}

$$s'_{k} = \frac{1}{p_{k}} \int_{u_{k}}^{u_{k+1}} s f(s) ds = \frac{1}{p_{k}} \cdot \frac{p_{k}}{u_{k+1} - u_{k}} \int_{u_{k}}^{u_{k+1}} s ds$$

$$= \frac{1}{2} \cdot \frac{1}{u_{k+1} - u_{k}} \cdot \left(u_{k+1}^{2} - u_{k}^{2}\right) = \frac{1}{2} \cdot \frac{\left(u_{k+1} + u_{k}\right) \cdot \left(u_{k+1} - u_{k}\right)}{u_{k+1} - u_{k}}$$

$$= \frac{1}{2} \left(u_{k} + u_{k+1}\right) = u_{k} + \frac{\Delta_{k}}{2}$$

$$D = \sum_{\forall k} \int_{u_k}^{u_{k+1}} (s - s'_k)^2 \cdot f(s) \, \mathrm{d}s$$

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$$D = \frac{1}{12} \sum_{\forall k} p_k \, \Delta_k^2$$

■ Will use: **Hölders inequality** in the following form (with $x_k \ge 0$ and $y_k \ge 0$)

$$\alpha + \beta = 1$$
 \Longrightarrow $\left(\sum_{k} x_{k}\right)^{\alpha} \cdot \left(\sum_{k} y_{k}\right)^{\beta} \geq \sum_{k} x_{k}^{\alpha} y_{k}^{\beta}$

with equality iff y_k is proportional to x_k , i.e., $y_k = \text{const} \cdot x_k$

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Average MSE distortion of Lloyd quantizer of size K (at high rates)

Approximation for centroid quantizers

$$D = \frac{1}{12} \sum_{i=0}^{K-1} p_i \, \Delta_i^2 = \frac{1}{12} \sum_{i=0}^{K-1} f(s_i') \, \Delta_i^3$$

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■ Rewrite expression using $\sum_{i=0}^{K-1} (1/K) = K \cdot (1/K) = 1$

$$D = \frac{1}{12} \left(\left(\sum_{i=0}^{K-1} f(s_i') \, \Delta_i^3 \right)^{\frac{1}{3}} \cdot \left(\sum_{i=0}^{K-1} \frac{1}{K} \right)^{\frac{2}{3}} \right)^3$$

Average MSE distortion of Lloyd quantizer of size K (at high rates)

Apply Hölders inequality

$$D \geq \frac{1}{12} \left(\sum_{i=0}^{K-1} \left(f(s_i') \, \Delta_i^3 \right)^{\frac{1}{3}} \left(\frac{1}{K} \right)^{\frac{2}{3}} \right)^3 = \frac{1}{12 \, K^2} \left(\sum_{i=0}^{K-1} \sqrt[3]{f(s_i')} \, \Delta_i \right)^3$$

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Average MSE distortion of Lloyd quantizer of size K (at high rates)

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■ Remember: Lloyd quantizer minimizes distortion for given size *K*

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■ Remember: Lloyd quantizer minimizes distortion for given size *K*

$$D = \frac{1}{12 \, \mathsf{K}^2} \left(\sum_{i=0}^{\mathsf{K}-1} \sqrt[3]{f(s_i')} \, \Delta_i \right)^3$$

■ Asymptotic limit for large K ($\Delta_k \rightarrow 0$)

$$D = \frac{1}{12 \, K^2} \left(\int_{-\infty}^{\infty} \sqrt[3]{f(s)} \, \mathrm{d}s \right)^3$$

High-Rate Approximation: Lloyd Quantizer with Fixed-Length Coding

MSE Distortion for Lloyd Quantizer at High Rates and Rate for Fixed-Length Coding

$$D = \frac{1}{12 \, K^2} \left(\int_{-\infty}^{\infty} \sqrt[3]{f(s)} \, \mathrm{d}s \right)^3 \qquad \text{and} \qquad R = \log_2 K \quad \Longrightarrow \quad \frac{1}{K^2} = 2^{-2R}$$

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Lloyd Quantizer with Fixed-Length Coding at High Rates

■ Panter and Dite approximation for operational distortion-rate function

$$D_F(R) = \frac{1}{12} \left(\int_{-\infty}^{\infty} \sqrt[3]{f(s)} \, \mathrm{d}s \right)^3 \cdot 2^{-2R}$$

$$D_F(R) = \varepsilon_F^2 \cdot \sigma^2 \cdot 2^{-2R}$$
 with

$$\varepsilon_F^2 = \frac{1}{12\,\sigma^2} \left(\int_{-\infty}^{\infty} \sqrt[3]{f(s)} \, \mathrm{d}s \right)^3$$

Lloyd Quantizer with Fixed-Length Coding vs Panter-Dite Approximation

Uniform pdf

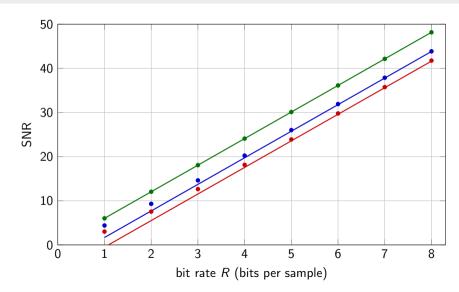
$$\varepsilon_F^2 = 1$$

Gaussian pdf

$$\varepsilon_F^2 = \frac{\sqrt{3}\,\pi}{2}$$

Laplacian pdf

$$\varepsilon_F^2 = 4.5$$



Summary of Lecture

Scalar Quantization

- Input-output function s' = Q(s) is a staircase function
- lacktriangle Quantizer is characterized by K reconstruction levels s_k' and K-1 decision thresholds u_k

Lloyd Quantizer

- Minimizes distortion *D* for given number *K* of quantization intervals
- Two optimization criterions
 - Centroid condition (MSE): $s'_k = \mathbb{E}\{S \mid S \in \mathcal{I}_k\}$
 - Nearest neighbor condition (MSE): $u_k = (s'_k + s'_{k-1})/2$
- Lloyd quantizer design: Iterate between the two optimization criterions
- High-rate approximation of Lloyd quantizer with fixed-length coding (Panter-Dite approximation)

Next Steps

- Theoretical limits for lossy source coding
- Consider entropy coding in quantizer design

Exercise 1: Implement Lloyd Algorithm

Implement the Lloyd algorithm using a programming language of your choice.

- Test the algorithm (for quantizer sizes of K = 2, 4, 8, 16, 32) for
 - a unit-variance Gaussian pdf:

$$f(s) = \frac{1}{\sqrt{2\pi}} e^{-\frac{1}{2} s^2}$$

• a unit-variance Laplacian pdf:

$$f(s) = \frac{1}{\sqrt{2}} \, e^{-\sqrt{2}\,|s|}$$

- Determine the distortion *D* for your quantizers.
- Compare the R-D performance of your quantizers (for K = 2, 4, 8, 16, 32) to the high-rate approximation for Lloyd quantizers with fixed-length codes.

You can implement the Lloyd algorithm that directly uses the pdf or the Lloyd algorithm that uses a training set (files with 1 000 000 samples in float32 format are provided on the course web site)

Exercise 2: Lloyd Quantizer for MSE Distortion (Alternative)

Given is a stationary source with a zero-mean Laplace pdf f(x) and a symmetric 3-interval quantizer:

$$f(x) = \frac{1}{\sqrt{2\sigma^2}} e^{-\sqrt{\frac{2}{\sigma^2}}|x|} \qquad \text{and} \qquad Q(x) = \begin{cases} -b : x < -a \\ 0 : |x| \le a \\ b : x > a \end{cases}$$

- (a) Derive the optimal reconstruction value b as a function of the threshold a for MSE distortion. Express the resulting distortion as function of the threshold a and the variance σ^2 .
- (b) Determine the decision threshold a in a way that a Lloyd quantizer for MSE distortion is obtained. Determine the distortion and rate for the Lloyd quantizer by assuming fixed-length coding $(R = \log_2 K)$ and compare the obtained R-D point with the high-rate approximation.
- (c) Can the derived optimal quantizer for fixed-length coding be improved by adding entropy coding (without changing the decision thresholds and reconstruction levels)?

Exercise 3: Lloyd Quantizer for MAE Distortion (Another Alternative)

Given is a stationary source with a zero-mean Laplace pdf f(x) and a symmetric 3-interval quantizer:

$$f(x) = \frac{1}{2m} e^{-\frac{|x|}{m}}$$
 and $Q(x) = \begin{cases} -b : x < -a \\ 0 : |x| \le a \\ b : x > a \end{cases}$

(a) Derive the centroid condition and nearest neighbor condition for MAE distortion

$$D = \mathrm{E}\{\,|S - S'|\,\}$$

- (b) Derive the optimal reconstruction value b as a function of the threshold a for MAE distortion. Express the resulting distortion as function of the threshold a and the parameter m.
- (c) Determine the decision threshold a in a way that a Lloyd quantizer for MAE distortion is obtained. Determine the distortion and rate for the quantizer by assuming fixed-length coding ($R = \log_2 K$).