Vector Quantization



Last Lectures: Scalar Quantization



■ Performance of Scalar Quantizers: Distortion (MSE) and Bit Rate

$$D = E\left\{\left(S - Q(S)\right)^{2}\right\} = \sum_{\forall k} \int_{u_{k}}^{u_{k+1}} (s - s_{k}')^{2} f(s) ds$$
$$R = E\left\{\ell\left(Q(S)\right)\right\} = \sum_{\forall k} \ell_{k} \int_{u_{k}}^{u_{k+1}} f(s) ds$$

Last Lectures: Optimal Scalar Quantization

Lloyd Quantizer

- Minimizes distortion D for given number K of reconstruction levels
- Two optimization criterions:
 - Centroid condition (MSE): $s'_k = \mathrm{E}\{ \ S \mid S \in \mathcal{I}_k \}$
 - Nearest neighbor condition (MSE): $u_k = (s'_k + s'_{k-1})/2$
- Lloyd quantizer design: Iterate between the two optimization criterions

Entropy-Constrained Lloyd Quantizer

- Minimizes rate-distortion cost $J = D + \lambda R$ for given Lagrange multiplier $\lambda > 0$
- Three optimization criterions:
 - Centroid condition (MSE): $s'_k = \mathrm{E}\{ \ S \mid S \in \mathcal{I}_k \}$
 - Entropy condition:

- $s'_{k} = \mathrm{E}\{ S \mid S \in \mathcal{I}_{k} \}$ $\ell_{k} = -\log_{2} \int_{u}^{u_{k+1}} f(s) \, \mathrm{d}s$
- Mod. nearest neighbor condition (MSE): $u_k = (s'_k + s'_{k-1})/2 + (\lambda/2)(\ell_k \ell_{k-1})/(s'_k + s'_{k-1})$
- EC-Lloyd quantizer design: Iterate between the optimization criterions

Last Lectures: Performance of Scalar Quantizers

High-Rate Approximations (MSE Distortion)

General form of high-rate distortion-rate function

$$D_X(R) = \varepsilon_X^2 \cdot \sigma^2 \cdot 2^{-2R}$$

where the constant factor ε_X^2 depends on shape of pdf and quantizer design

- → Lloyd + fixed length: $\varepsilon_F^2 = \frac{1}{12} \left(\int_{-\infty}^{\infty} \sqrt[3]{f(s/\sigma)} \, ds \right)^3$
- → EC-Lloyd + variable length: $\varepsilon_V^2 = \frac{1}{12} 2^{2h(S/\sigma)}$ with $h(S) = -\int_{-\infty}^{\infty} f(s) \log_2 f(s) ds$
- → Shannon lower bound: $\varepsilon_L^2 = \frac{1}{2\pi e} 2^{2 h(S/\sigma)}$

Comparison of Coding Efficiency

- EC-Lloyd often significantly better than Lloyd (Gauss: 2.81 dB; Laplace: 5.63 dB)
- Constant high-rate performance gap between EC-Lloyd and Shannon lower bound

$$\frac{D_V(R)}{D_L(R)} = \frac{\pi e}{6} \approx 1.42 \ (1.53 \text{ dB}), \qquad \qquad R_V(D) - R_L(D) = \frac{1}{2} \log_2 \frac{\pi e}{6} \approx 0.25$$

Last Lectures: Scalar Quantization in Practice



Uniform Reconstruction Quantizers (URQs)

- Simple decoding process: $s'_n = \Delta \cdot q_n$
- Encoder can choose trade-off between coding efficiency and complexity
 - → Simplest encoding: $q_n = \text{round} (s_n / \Delta)$
 - → Optimal encoding: Choose q_n that minimizes Lagrange cost $J(q_n) = (s_n q_n \cdot \Delta)^2 + \lambda \cdot \ell_k$, typically using fixed relationship $\lambda = \text{const} \cdot \Delta^2$

URQs with optimal encoding are virtually as good as optimal scalar quantizers (for typical pdfs)

Quantization: Open Questions

Performance Gap to Theoretical Bound

Remember: High-rate performance of optimal scalar quantizer for IID sources

$$rac{D_V}{D_L}\left(R
ight)=rac{\pi e}{6}pprox 1.42 \qquad (1.53\,\mathrm{dB}\,\log s)$$
 (1.53 (B)

- → What causes this performance gap?
- → How can the quantizer performance be improved?

Quantization of Sources with Memory

- Scalar quantizers cannot exploit dependencies between samples (use only marginal pdf)
- → How can we improve lossy coding for sources with memory?
 - Conditional entropy coding of quantization indexes?
 - Combination of scalar quantization and prediction?

• ... ?

Scalar Quantizers in N-dimensional Signal Space



Interpretation of Scalar Quantization in N-dimensional Signal Space

- N-dimensional input vector s is mapped to N-dimensional reconstructed vector s'
- All vectors s inside a quantization cell C_k are mapped to the same reconstruction vector s'_k
- → Quantization cells C_k form hyper-rectangles in N-dimensional signal space
- \rightarrow Reconstruction vectors s'_k lie on orthogonal grid aligned with coordinate axes

Vector Quantization: Relaxing Structural Constraints



Vector Quantization

- Joint quantization of vectors/blocks s of N > 1 successive input samples
- Relax structural constraints that are implicitly imposed by scalar quantization
 - \bullet Quantization cells \mathcal{C}_k can be arbitrarily shaped in N-dimensional space
 - Reconstruction vectors s'_k can be arbitrarily placed in N-dimensional space

→ Allows a number of new options in designing quantizers

Structure of Vector Quantizers

Vector Quantizers of Quantizer Dimension N

• Map N-d input vectors s to N-d output vectors s'_k

$$Q: \mathbb{R}^N \mapsto \{ \mathbf{s_0'}, \mathbf{s_1'}, \mathbf{s_2'}, \cdots \}$$

Partition N-d space into countable number of quantization cells C_k

$$\mathcal{C}_k = \{ \, \boldsymbol{s} \in \mathbb{R}^N : \ \boldsymbol{Q}(\boldsymbol{s}) = \boldsymbol{s}'_{\boldsymbol{k}} \, \}$$

All input vectors s that fall inside a quantization cell C_k are mapped to the associated reconstruction vector s'_k

Vector Quantization and Entropy Coding

- Quantization index k indicates quantization cell C_k and reconstruction vector s'_k
 - → Encoder mapping: $\alpha(s) = k$, $\forall s \in C_k$
 - → Decoder mapping: $\beta(k) = s'_k$



Vector Quantization: Encoding and Decoding

- Arbitrarily shaped quantization cells C_k are difficult to store and check
- → Concept of quantization cells is not required in practice

Encoding

- Select the reconstruction vector s' that minimizes a distance measure d to the input vector s
- Possible distance measures:
 - → MSE distortion: $d = ||\boldsymbol{s} \boldsymbol{s'_k}||_2^2$
 - → Lagrangan cost: $d = ||\boldsymbol{s} \boldsymbol{s'_k}||_2^2 + \lambda \ell_k$

Decoding

 Output reconstruction vector s'_k indicated by transmitted quantization index k (use array in decoder)



Vector Quantization / Performance of Vector Quantizers

Performance of Vector Quantizers: Bit Rate

- Let ℓ_k be the codeword length for quantization index k
- → Average bit rate *R* per sample

$$R = \frac{1}{N} \operatorname{E} \left\{ \ell(Q(\boldsymbol{S})) \right\} = \frac{1}{N} \sum_{\forall k} p_k \ell_k$$

Probability p_k of quantization cell C_k / quant. index k

$$p_k = \int_{\mathcal{C}_k} f(\boldsymbol{s}) \, \mathrm{d} \boldsymbol{s}$$

Approximation for training set

$$p_k = \frac{n(k)}{\sum_k n(k)}$$

where n(k) is the number of vectors assigned to C_k / s'_k



Performance of Vector Quantizers: Distortion

→ Average MSE distortion *D* per sample

$$D = \frac{1}{N} \operatorname{E} \left\{ \left| \left| \mathbf{S} - Q(\mathbf{S}) \right| \right|_{2}^{2} \right\}$$
$$= \frac{1}{N} \int_{\mathbb{R}^{N}} \left| \left| \mathbf{s} - Q(\mathbf{s}) \right| \right|_{2}^{2} f(\mathbf{s}) \, \mathrm{d}\mathbf{s}$$
$$= \frac{1}{N} \sum_{\forall k} \int_{\mathcal{C}_{k}} \left| \left| \mathbf{s} - \mathbf{s}_{k}^{\prime} \right| \right|_{2}^{2} f(\mathbf{s}) \, \mathrm{d}\mathbf{s}$$
$$D = \frac{1}{N} \sum_{\forall k} \int_{\mathcal{C}_{k}} \left(\mathbf{s} - \mathbf{s}_{k}^{\prime} \right)^{\mathrm{T}} \left(\mathbf{s} - \mathbf{s}_{k}^{\prime} \right) f(\mathbf{s}) \, \mathrm{d}\mathbf{s}$$

• Approximation for training set $\{s_n\}$ of L vectors

$$D = rac{1}{L} \sum_{orall n} \left| \left| oldsymbol{s}_n - Q(oldsymbol{s}_n) \right|
ight|_2^2$$



Optimal Vector Quantizer for Fixed-Length Coding

Goal: Minimize MSE Distortion for K Quantization Cells

- Similar to Scalar Lloyd Quantizer
- Neglect impact of entropy coding → Consider fixed-length coding
- \rightarrow Rate *R* and MSE distortion *D* are given by

$$R = \frac{1}{N} \left[\log_2 K \right] \qquad \text{(typically } K = 2^B \text{, with } B \text{ being the bits per codeword)}$$
$$D = \frac{1}{N} \sum_{\forall k} \int_{\mathcal{C}_k} \left| \left| s - s'_k \right| \right|_2^2 f(s) \, \mathrm{d}s = \frac{1}{N} \sum_{\forall k} \int_{\mathcal{C}_k} \left(s - s'_k \right)^{\mathrm{T}} \left(s - s'_k \right) f(s) \, \mathrm{d}s$$

Optimize Quantizer of size K

- Derive necessary conditions for optimality (similar to Lloyd quantizer)
- Construct iterative algorithm for designing quantizer

Optimality Conditions for Fixed-Length Coding

Necessary Conditions for Optimality (MSE distortion)

1 Centroid condition (for reconstruction vectors s'_k)

$$oldsymbol{s_k'} = \mathrm{E}\{ egin{array}{c} oldsymbol{s} \mid oldsymbol{S} \in \mathcal{C}_k \ \} = rac{1}{p_k} \int_{\mathcal{C}_k} oldsymbol{s} \, f(oldsymbol{s}) \, \mathrm{d}oldsymbol{s} \end{array}$$

➡ Centroid condition for training set

$$m{s}_{m{k}}' = rac{1}{n(k)} \sum_{orall s: \ lpha(m{s})=k} m{s} \qquad ext{with} \qquad n(k) = \sum_{orall s: \ lpha(m{s})=k} m{1}$$

2 Nearest neighbour condition (for quantization cells C_k / encoder mapping $\alpha(.)$)

$$lpha(\boldsymbol{s}) = \arg\min_{orall k} ||\boldsymbol{s} - \boldsymbol{s'_k}||_2^2$$

The Linde-Buzo-Gray (LBG) Algorithm for a Training Set (MSE Distortion)

Given is: • the dimension N and the size K of the quantizer

• a sufficiently large realization $\{s_n\}$ of considered source

Iterative quantizer design

- **1** Choose an initial set of K reconstruction vectors $\{s'_k\}$
- **2** Associate all vectors of the training set $\{s_n\}$ with one of the quantization cells C_k

$$q(\boldsymbol{s}_n) = \arg\min_{orall k} \left|\left|\boldsymbol{s}_n - \boldsymbol{s}'_{\boldsymbol{k}}\right|\right|_2^2$$
 (nearest neighbor condition)

3 Update the reconstruction vectors $\{s'_k\}$ according to

$$m{s}_{m{k}}' = rac{1}{n(k)} \sum_{\forall n: \ m{q}(m{s}_n) = k} m{s}_n$$
 (centroid condition)

where n(k) is the number of sample vectors \boldsymbol{s}_n assigned to C_k

4 Repeat the previous two steps until convergence

Example: LBG Algorithm for Gaussian IID ($\sigma^2 = 1$, N = 2, K = 16)



Comparison to Scalar Quantization: Gaussian IID ($\sigma^2 = 1$, R = 2)



→ Improvement of 0.39 dB (distortion reduction by factor 0.91)

Comparison to Scalar Quantization: Gaussian IID ($\sigma^2 = 1$, R = 4)



LBG (N = 2)-2 _4 -20 2 -4D = 0.00767 $SNR = 21.15 \, dB$

→ Improvement of 0.93 dB (distortion reduction by factor 0.81)

Comparison to Scalar Quantization: Laplacian IID ($\sigma^2 = 1$, R = 2)



→ Improvement of 1.35 dB (distortion reduction by factor 0.73)

Comparison to Scalar Quantization: Laplacian IID ($\sigma^2 = 1, R = 4$)



LBG (N = 2)



→ Improvement of 1.94 dB (distortion reduction by factor 0.64)

The Vector Quantizer Advantage

Gain over scalar quantization can be assigned to 3 effects:

Space filling advantage:

- \mathbb{Z}^N lattice is not most efficient sphere packing in N dimensions (N > 1)
- Independent from source distribution or statistical dependencies
- Maximum gain for $N \to \infty$: 1.53 dB

Shape advantage:

- Exploit shape of source pdf
- Can also be exploited using entropy-constrained scalar quantization

Memory advantage:

- Exploit statistical dependencies of the source
- Can also be exploited using predictive coding, transform coding, block entropy coding or conditional entropy coding

Space-Filling Advantage: LBG for Uniform IID Source

Lloyd (N = 1): SNR = 23.97 dB



LBG (N = 2): SNR = 24.14 dB



- LBG algorithm approaches approximate hexagonal lattice
- → Improvement of 0.17 dB

Space-Filling Advantage: Sphere Packing in N-dimensional Signal Space





- Space filling gain: Densest packing of "optimal" quantization cells in signal space
- → MSE distortion: Densest packing of spheres in *N*-dimensional space
 - → 2 dimensions: Hexagonal lattice (like honeycombs)
 - → 3 dimensions: Cuboidal lattice (stapling of cannon balls / oranges)

Space-Filling Advantage: Sphere Packing Density

Center density

- Consider *N*-dimensional spheres with radius r = 1
- Measure for packing density: Center density

 $\delta = \frac{\text{average number of sphere centers}}{\text{unit volume}}$

• Example:
$$N = 1$$
 (SQ with intervals of size $2r = 2$)

$$\delta = \frac{1}{2}$$

Roger's bound

• Theoretical upper bound for center density (last term being approximate)

$$\log_2 \delta \leq \frac{N}{2} \log_2 \left(\frac{N}{4e\pi} \right) + \frac{1}{2} \log_2 \left(\frac{\pi N^3}{e^2} \right) + \frac{21}{4N + 10}$$

Space-Filling Advantage: Densest Known Sphere Packings

- Densest known packings for dimensions $N \le 48$ [Conway, Sloane, 1998]
- Vertical axis: $\log_2 \delta + N(24 N)/96$



Space-Filling Advantage: Approximate SNR Gain

dimension	densest packing	(name)	highest kissing number	approximate gain [dB]
1	Z	– Integer lattice	2	0
2	A_2	– Hexagonal lattice	6	0.17
3	$A_3 \simeq D_3$	– Cuboidal lattice	12	0.29
4	D_4		24	0.39
5	D_5		40	0.47
6	E ₆		72	0.54
7	E ₇		126	0.60
8	E ₈	- Gosset lattice	240	0.66
9	Λ_9	 Laminated lattice 	240	0.70
10	P _{10c}	 Non-lattice arrangement 	336	0.74
12	K ₁₂	 Coxeter-Todd lattice 	756	0.81
16	$BW_{16}\simeq \Lambda_{16}$	– Barnes-Wall lattice	4320	0.91
24	Λ_{24}	- Leech lattice	196560	1.04
100				1.35
∞				1.53

Summary on Space-Filling Advantage

- Gain of unique to vector quantization: Packing of quantization cells in *N*-dimensional space
- Increases with quantizer dimension N
- → Gain for $N \rightarrow \infty$: Difference between Shannon lower bound and ECSQ



Shape Advantage: Gaussian IID ($\sigma^2 = 1, R = 4$)



Shape Advantage of Vector Quantizers

- Coding gain (0.93 dB for example) is larger than space-filling gain (0.17 dB for N = 2)
- Vector quantizer can better adapt to shape of pdf (even without entropy coding)

Shape Advantage: Laplacian IID ($\sigma^2 = 1, R = 4$)



Shape Advantage of Vector Quantizers

- Coding gain (1.94 dB for example) is larger than space-filling gain (0.17 dB for N = 2)
- Vector quantizer can better adapt to shape of pdf (even without entropy coding)

Summary on Shape Advantage

- Gain of VQ due to exploitation of shape of pdf (without entropy coding)
- Overall gain for iid source: Space-filling gain + shape gain
- → Shape advantage can also be exploited by entropy-constrained scalar quantization



Memory Advantage: Gauss-Markov ($\sigma^2 = 1$, $\varrho = 0.9$, R = 2)



Memory Advantage of Vector Quantizers

- Large coding gain (4.20 dB for example) for sources with memory
- Vector quantizer can exploit dependencies between samples

Memory Advantage: Gauss-Markov ($\sigma^2 = 1$, $\varrho = 0.9$, R = 4)



Memory Advantage of Vector Quantizers

- Large coding gain (4.55 dB for example) for sources with memory
- Vector quantizer can exploit dependencies between samples

Summary on Memory Advantage

- Gain of VQ due to exploitation of dependencies between samples
- Largest gain to be made for sources with strong statistical dependencies
- → Exploitation of memory advantage is one of the most relevant aspects in source coding



Optimal Vector Quantizer for Variable-Length Coding

Optimal Vector Quantizer with Consideration of Entropy Coding

- Similar to Scalar Entropy-Constrained Lloyd Quantizer
- ightarrow Minimization of Lagrangian cost for given Lagrange multiplier λ

$$J = D + \lambda \cdot R$$

= $\frac{1}{N} \sum_{\forall k} \int_{\mathcal{C}_k} ||\boldsymbol{s} - \boldsymbol{s}'_k||_2^2 f(\boldsymbol{s}) \, \mathrm{d}\boldsymbol{s} + \frac{\lambda}{N} \sum_{\forall k} \ell_k \int_{\mathcal{C}_k} f(\boldsymbol{s}) \, \mathrm{d}\boldsymbol{s}$

• Lagrange multiplier $\lambda > 0$ determines operation point (trade-off between rate and distortion)

Optimize Quantizer for given Lagrange multiplier

- Derive necessary conditions for optimality (similar to EC Lloyd quantizer)
- Construct iterative algorithm for designing quantizer
- Similar as for EC Lloyd: Use large number of intervals in initialization

Optimality Conditions for Variable-Length Coding

Necessary Conditions for Optimality (MSE distortion)

1 Centroid condition (for reconstruction vectors s'_k , same as for LBG)

$$oldsymbol{s_k'} = \mathrm{E}\{ egin{array}{c} oldsymbol{s} \mid oldsymbol{S} \in \mathcal{C}_k \ \} = rac{1}{p_k} \int_{\mathcal{C}_k} oldsymbol{s} \, f(oldsymbol{s}) \, \mathrm{d}oldsymbol{s} \end{array}$$

(training set: take average of assigned vectors)

2 Entropy condition (for codeword length ℓ_k , same as for EC Llloyd)

$$\ell_k = -\log_2 p_k = -\log_2 \int_{\mathcal{C}_k} f(\boldsymbol{s}) \, \mathrm{d} \boldsymbol{s}$$

(training set: count assigned vectors)

3 Modified nearest neighbour condition (for quantization cells C_k / encoder mapping $\alpha(.)$)

$$lpha(oldsymbol{s}) = rg \min_{orall k} \left| \left| oldsymbol{s} - oldsymbol{s}'_{oldsymbol{k}}
ight| \right|_2^2 + \lambda \cdot \ell_k
ight|$$

The Chou-Lookabough-Gray (CLG) Algorithm for a Training Set (MSE)

- Given is: the Lagrange multiplier $\lambda > 0$
 - a sufficiently large realization $\{s_n\}$ of considered source

Iterative quantizer design

- **1** Choose an initial set of K reconstruction vectors $\{s'_k\}$ and codeword length $\{\ell_k\}$
- **2** Associate all vectors of the training set $\{s_n\}$ with one of the quantization cells C_k

$$q(\boldsymbol{s}_n) = \arg\min_{\forall k} \left| \left| \boldsymbol{s}_n - \boldsymbol{s}'_k \right| \right|_2^2 + \lambda \cdot \ell_k \qquad (\text{modified nearest neighbor condition})$$

3 Update the reconstruction vectors $\{s'_k\}$ and codeword length $\{\ell_k\}$ according to

$$m{s'_k} = rac{1}{n(k)} \sum_{orall n: \ q(m{s}_n) = k} m{s}_n \qquad ext{ and } \qquad \ell_k = -\log_2\left(rac{n(k)}{\sum_{orall i} n(i)}
ight)$$

where n(k) is the number of sample vectors s_n assigned to C_k

4 Repeat the previous two steps until convergence

Entropy-Constrained Vector Quantizer: Gaussian IID ($\sigma^2 = 1$, R = 2)



→ Large gain (1.0 dB) relative to LBG algorithm (fixed-length coding)

→ Gain relative to EC Lloyd reduces to space-filling gain (0.17 dB for N = 2)

Entropy-Constrained Vector Quantizer: Gaussian IID ($\sigma^2 = 1$, R = 3)



 \rightarrow Large gain (1.4 dB) relative to LBG algorithm (fixed-length coding)

→ Gain relative to EC Lloyd reduces to space-filling gain (0.17 dB for N = 2)

Entropy-Constrained Vector Quantizer: Gauss-Markov ($\sigma^2 = 1$, $\varrho = 0.9$, R = 3)



 \Rightarrow Large gain (1.0 dB) relative to LBG algorithm (fixed-length coding)

→ Gain relative to EC Lloyd: Sum of memory gain and space-filling gain

Coding Efficiency of Vector Quantizers

Vector Quantizer Advantages

- Space-filling advantage
 - → Unique to vector quantization: (0.17 dB for N = 2; 1.53 dB for $N \rightarrow \infty$)
- Shape advantage
 - → Can also be exploited by entropy-constrained scalar quantization
- Memory advantage
 - → Can also be (partly) exploited by other coding techniques (topic of next lectures)

Coding Efficiency of Vector Quantizers

- Optimal vector quantizers provide coding efficiency gains relative to scalar quantizers
- IID sources: Only space-filling gain (when comparing entropy-constrained designs)
- Sources with memory: Most important aspect is the memory advantage

ightarrow Vector quantizers can asymptotically achieve rate-distortion bound for $N ightarrow\infty$

Complexity of Vector Quantization

Decoding Complexity

- In principle: Table look-up (using transmitted quantization indexes)
- Extremely large memory requirements for large N

Encoding Complexity

- Finding the "closest reconstruction vector" can become very complex
- Designing a good vector quantizer is already very complex

Usage of Vector Quantization

- Unconstrained vector quantizers are rarely used in practice
- → Reduce complexity by imposing structural constraints
 - Tree-structured vector quantizers
 - Gain-shape vector quantizers
 - Lattice vector quantizers (important special case: Transform coding)
 - Trellis-coded quantization

Lattice Vector Quantizers & Transform Coding

Lattice Vector Quantizer

- Reconstruction vectors are located on multi-dimensional lattice
 - Lattice is specified by N "basis vectors" {**b**_k}
 - Reconstruction vectors given by matrix of "basis vectors"

$$\boldsymbol{s}'_{\boldsymbol{k_1},\boldsymbol{k_2},\cdots,\boldsymbol{k_N}} = \boldsymbol{M} \cdot \left[k_1,k_2,\cdots,k_N\right]^{\mathrm{T}}$$

- Simple decoder operation possible
- Less complex encoding (can still by very complex for large N)

Transform Coding

- Lattice vector quantizer with orthonormal "basis vectors"
- Very simple encoding and decoding
- ightarrow One of the most often used approaches in lossy coding
- \rightarrow Will discuss in detail in next lectures



Trellis-Coded Quantization (TCQ)

Quantizer Design & Decoding Process

Two scalar quantizers + Prodecure for switching between quantizers (state machine with 2^N states)



state	quantizer	next state
0	Q_0	$(A,B)\mapsto (0,1)$
1	Q_1	$(C,D)\mapsto(2,3)$
2	Q_0	$(A,B)\mapsto(1,0)$
3	Q_1	$(C,D)\mapsto (3,2)$

Encoding Process

■ Trellis formulation of possible quantizer switching → Viterbi algorithm



Example: TCQ Performance for Gaussian IID



Summary of Lecture

Vector Quantization (VQ)

- Straightforward extension of scalar quantization to higher dimensions N
- Opt. VQ with fixed-length codes: Similar to Lloyd quantizer
- Opt. VQ with variable-length codes: Similar to EC-Lloyd quantizer

Vector Quantizer Advantages

- Space-filling advantage: Unique to vector quantizers (1.53 dB for $N \to \infty$)
- Shape advantage: Can also be exploited by ECSQ
- Memory advantage: Can also be exploited by other coding techniques

Vector Quantization can achieve Rate-Distortion Bound! - Are we done?

- → No! Complexity of vector quanzization is a serious issue!
- ➡ Require lossy coding techniques with high rate-distortion efficiency and a complexity suitable for wide range of implementations
- ➡ Particularly important: Exploitation of dependencies between samples!

Exercise 1: Space-Filling Gain for 2-dimensional Vector Quantizer

Calculate the gain (in signal-to-noise ratio) of optimal 2-dimensional vector quantization relative to optimal scalar quantization for high rates on the example of a uniform pdf.

Hints:

- In two dimensions, the optimal quantization cells are regular hexagons; the associated reconstruction vectors are located in the centers of the hexagons.
- For high rates, border effects can be neglected. It can be assumed that the signal space for which the pdf is non-zero is completely filled with regular quantization cells.