

German University in Cairo - GUC Faculty of Information Engineering & Technology - IET Department of Communication Engineering Winter Semester 2013/2014

Students Name:

Students ID:

COMM901 – Source Coding and Compression

Quiz 2

Note:

- The quiz contains 4 questions (each of which has 10 points)
- Only the best 3 questions are counted (maximum grade is 30)

	Question 1	Question 2	Question 3	Question 4
possible points	10	10	10	10
achieved point				
counted (best 3Q)				
			sum counted:	
			percentage:	

<u>Question 1:</u> Mutual Information [10 points]

Given is a stationary Markov process $S = \{S_n\}$ with the binary symbol alphabet $\mathcal{A}_S = \{x, y\}$. The conditional symbol probabilities $p(s_n|s_{n-1}) = P(S_n = s_n|S_{n-1} = s_{n-1})$ are given in the left table below. Due to the symmetry of the transition probabilities, the marginal probability mass function is uniform; the corresponding marginal symbol probabilities $p(s_n) = P(S_n = s_n)$ are given in the right table below.

<i>S</i> _n	$p(s_n s_{n-1} = x)$	$p(s_n s_{n-1} = y)$	<i>S</i> _n	$p(s_n)$
x	3/4	1/4	x	1/2
у	1/4	3/4	у	1/2

Determine the mutual information $I(S_n; S_{n+1})$ between two successive random variables S_n and S_{n+1} of the given stationary Markov source.

The mutual information $I(S_n; S_{n+1})$ between two successive samples is given by

$$I(S_n; S_{n+1}) = H(S_{n+1}) - H(S_{n+1}|S_n)$$
 [2]

The marginal entropy $H(S_{n+1})$ is given by

$$H(S_{n+1}) = H(S_n) = -\sum_{\forall i} p(s_i) \cdot \log_2 p(s_i) = -2 \cdot \frac{1}{2} \cdot \log_2 \frac{1}{2} = 1$$
 [2]

For the conditional entropies $H(S_{n+1}|S_n = z)$, with z being x or y, we obtain

$$H(S_{n+1}|S_n = z) = -\sum_{\forall i} p(s_i|z) \cdot \log_2 p(s_i|z) = -\frac{3}{4} \cdot \log_2 \frac{3}{4} - \frac{1}{2} \cdot \log_2 \frac{1}{2}$$
$$= \frac{3}{2} - \frac{3}{4} \cdot \log_2 3 + \frac{1}{2} = 2 - \frac{3}{4} \cdot \log_2 3 \approx 0.8113$$
 [2]

Then, the conditional entropy $H(S_{n+1}|S_n)$ is given by

$$H(S_{n+1}|S_n) = \sum_{\forall i} p(a_i) \cdot H(S_{n+1}|S_n = a_i) = H(S_{n+1}|S_n = z) \cdot \sum_{\forall i} p(a_i)$$
$$= H(S_{n+1}|S_n = z) = 2 - \frac{3}{4} \cdot \log_2 3 \approx 0.8113$$
 [2]

Using the derived expression for the marginal entropy $H(S_n)$ and the conditional entropy $H(S_{n+1}|S_n)$, we obtain for the mutual information

$$I(S_n; S_{n+1}) = 1 - 2 + \frac{3}{4} \cdot \log_2 3 = \frac{3}{4} \cdot \log_2 3 - 1 \approx 0.1887$$
 [2]

The mutual information between two successive random variables is 0.1887 bit per sample.

Question 2: Properties of Rate-Distortion Functions [10 points]

Given are 5 random processes with independent and identically distributed random variables:

- The random process *A* is continuous and has a Laplace distribution.
- The random process **B** is continuous and has a Gaussian distribution.
- The random process *C* is discrete, has the alphabet $\{x, y\}$ and the pmf p(x) = p(y) = 1/2.
- The random process **D** is discrete, has the 8-symbol alphabet $\mathcal{A}_D = \{a, b, c, d, e, f, g, h\}$ and a uniform pmf on that alphabet (i.e., all symbols have the same probability).
- The random process *E* is discrete, has the infinite symbol alphabet $\mathcal{A}_E = \{0, 1, 2, 3, \dots\}$ and the geometric pmf $p(k) = (1/2)^{k+1}$, with $k \in \mathcal{A}_E$.

The following diagrams show sketches of the (information) distortion-rate function for four of the five given random processes. As distortion measure the mean squared error (MSE) is used. In the diagrams, the distortion D is represented as signal-to-noise ratio (SNR), given by

$$SNR = 10 \cdot \log_{10} \frac{\sigma^2}{D}$$

where σ^2 represents the variance of the random variables.

Assign to each of the diagrams the correct random process (*A*, *B*, *C*, *D*, or *E*).



Question 3: Usage of Rate-Distortion Functions [10 points]

A student of the German University of Cairo developed a lossy coding algorithm for arbitrary sources. His professor tested the algorithm for a long realization of a stationary Gauss-Markov process and measured a bit rate of R = 2 bit per sample and a signal-to-noise ratio of 20 dB.

Given this result, determine the possible values (or intervals) for the correlation coefficient ρ of the tested Gauss-Markov process.

<u>Hints:</u> For bit rates greater than 1 bit per sample, the information distortion-rate function for stationary Gauss-Markov processes and MSE distortion is given by $D(R) = (1 - \varrho^2) \cdot \sigma^2 \cdot 2^{-2R}$ The signal-to-noise ratio (SNR) is defined by $SNR = 10 \cdot \log_{10} \frac{\sigma^2}{MSE}$

The distortion-rate function D(R) specifies the minimum distortion that can be achieved by coding a source at a rate R. Hence, with $SNR_{D(R)}$ being the SNR associated with the distortion-rate function and SNR_X being the SNR of the actual coding result, we have

$$\text{SNR}_X \le \text{SNR}_{D(R)} = 10 \cdot \log_{10} \frac{\sigma^2}{D(R)} = -10 \cdot \log_{10} \left((1 - \varrho^2) \cdot 2^{-2R} \right)$$
 [4]

By inserting the given values for R and SNR_X , we obtain

$$20 \le -10 \cdot \log_{10} \left((1 - \varrho^2) \cdot 2^{-4} \right)$$
$$-2 \ge \log_{10} \left(\frac{(1 - \varrho^2)}{16} \right)$$
$$\frac{16}{100} \ge 1 - \varrho^2$$
$$\varrho^2 \ge 0.84 \quad [3]$$

This resulting quadratic inequality has the solutions $|\varrho| \ge \sqrt{0.84}$, which is equivalent to

$$\varrho \ge \sqrt{0.84} \approx 0.9165$$
 and $\varrho \le -\sqrt{0.84} \approx -0.9165$

Since a correlation coefficient is always inside the closed interval [-1; 1], it follows that, based on the coding results, the correlation coefficient ρ lies inside one of the following intervals:

$$-1 \le \varrho \le \sqrt{0.84}$$
 or $\sqrt{0.84} \le \varrho \le 1$ [3]

Or in other words, the absolute value of q is bounded by $\sqrt{0.84} \le |q| \le 1$.

Question 4: Shannon Lower Bound [10 points]

Given is a continuous-amplitude random process $X = \{X_n\}$ with independent and identically distributed random variables X_n . As illustrated in the figure below, the random variables X_n have a uniform probability density function in an interval [-a; a].



Determine the Shannon lower bound as function of the variance σ^2 . The Shannon lower bound should be formulated as distortion-rate function $D_L(R)$ with the MSE as distortion measure.

Hint: For MSE distortion, the Shannon lower bound as distortion-rate function is given by
$$D_L(R) = \frac{1}{2\pi e} \cdot 2^{2 \cdot \overline{h}(X)} \cdot 2^{-2R}$$

Let b be the value of the pdf inside the interval [-a; a]. Since the integral is equal to 1, we have

$$\int_{-a}^{a} f_X(x) \, \mathrm{d}x = b \cdot \int_{-a}^{a} \mathrm{d}x = 2ab = 1 \qquad \Longrightarrow \qquad b = \frac{1}{2a} \quad [\mathbf{2}]$$

Since we consider an iid process, the differential entropy rate is equal to the differential entropy:

$$\bar{h}(\mathbf{X}) = h(X) = -\int_{-a}^{a} f_{X}(x) \cdot \log_{2} f_{X}(x) \, dx = -\int_{-a}^{a} \frac{1}{2a} \cdot \log_{2} \frac{1}{2a} \, dx$$
$$= \frac{\log_{2}(2a)}{2a} \cdot \int_{-a}^{a} dx = \frac{\log_{2}(2a)}{2a} \cdot \left(a - (-a)\right) = \log_{2}(2a) \quad [3]$$

Hence, we obtain for the Shannon lower bound

$$D_L(R) = \frac{1}{2\pi e} \cdot 2^{2 \cdot \overline{h}(\mathbf{X})} \cdot 2^{-2R} = \frac{1}{2\pi e} \cdot 2^{2\log_2(2a)} \cdot 2^{-2R} = \frac{4a^2}{2\pi e} \cdot 2^{-2R} = \frac{2a^2}{\pi e} \cdot 2^{-2R}$$
[2]

Since we want to formulate the Shannon lower bound as function of the variance σ^2 , we calculate the variance σ^2 . Note that the mean is $\mu = 0$.

$$\sigma^{2} = \int_{-a}^{a} x^{2} \cdot f_{X}(x) \, \mathrm{d}x = \frac{1}{2a} \int_{-a}^{a} x^{2} \, \mathrm{d}x = \frac{1}{2a} \cdot \left(\frac{a^{3}}{3} - \left(-\frac{a^{3}}{3}\right)\right) = \frac{a^{2}}{3} \quad [2]$$

Hence, we have $a^2 = 3\sigma^2$, which yields the Shannon lower bound

$$D_L(R) = \frac{6}{\pi e} \cdot \sigma^2 \cdot 2^{-2R} \quad [1]$$